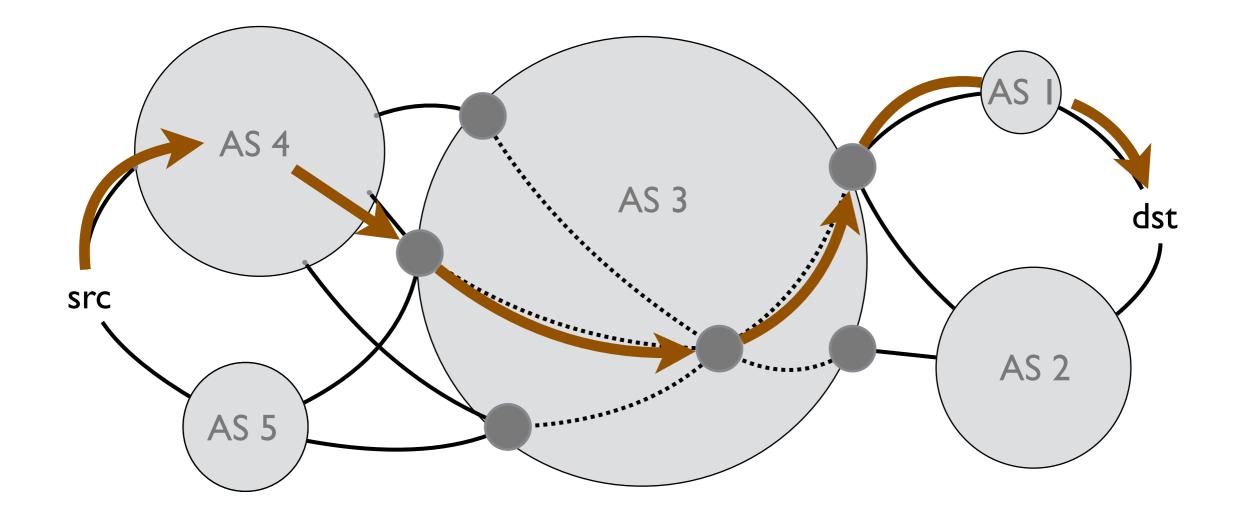
Enabling Policy Innovation in Interdomain Routing: A Software-Defined Approach

Anduo Wang^{*} Zhijia Chen^{*} Tony Yang[†] Minlan Yu[‡] *Temple University [†]Johns Hopkins University [‡]Harvard University

April 4, 2019 SOSR'19

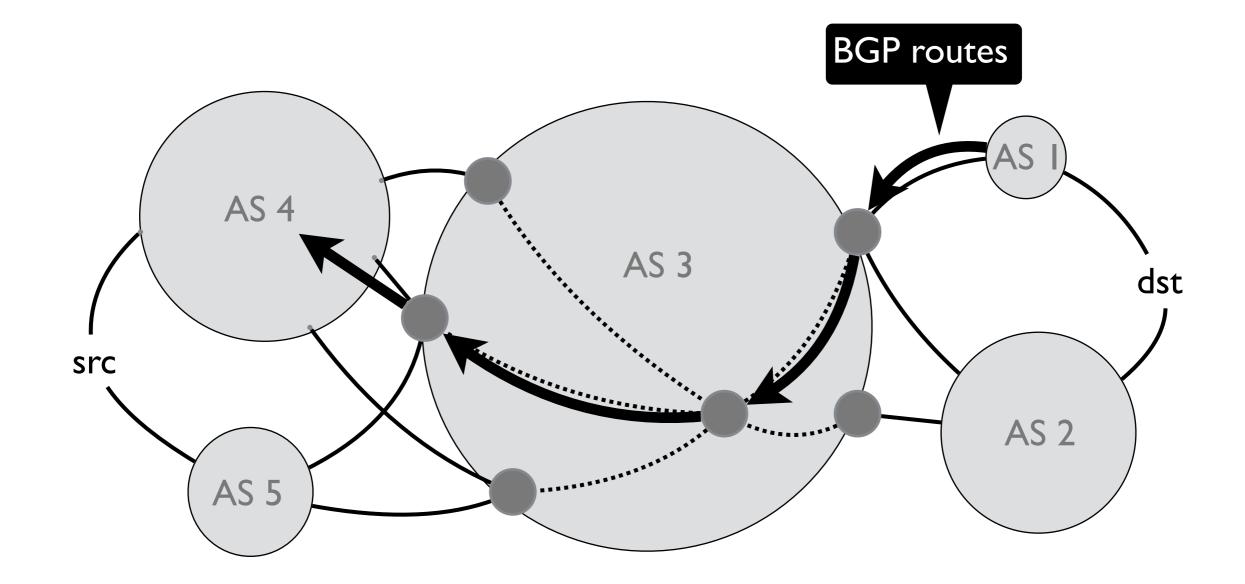
interdomain routing

determining data path connecting communicating hosts



interdomain routing

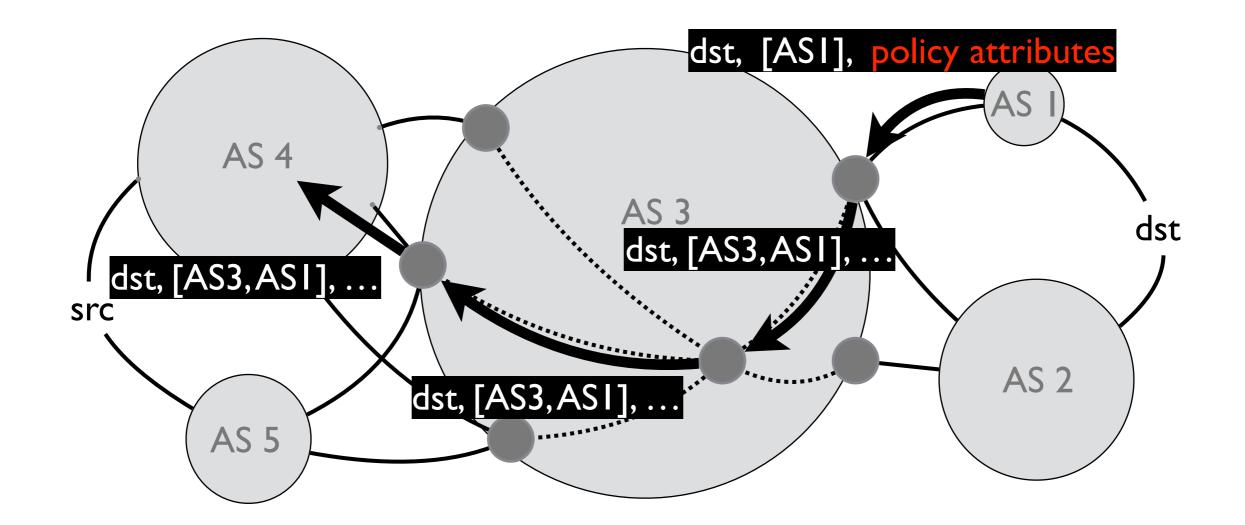
border gateway protocol (BGP) — the only defacto interdomain routing system



BGP and AS policy

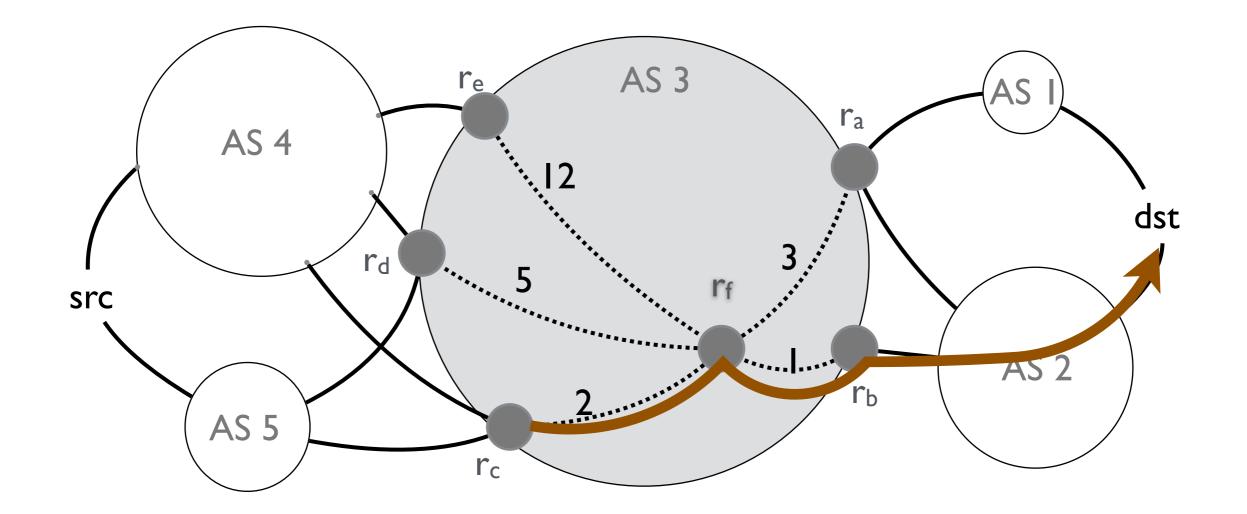
BGP supports AS policies

- extending path-vector system
- overriding the shortest AS-path behavior



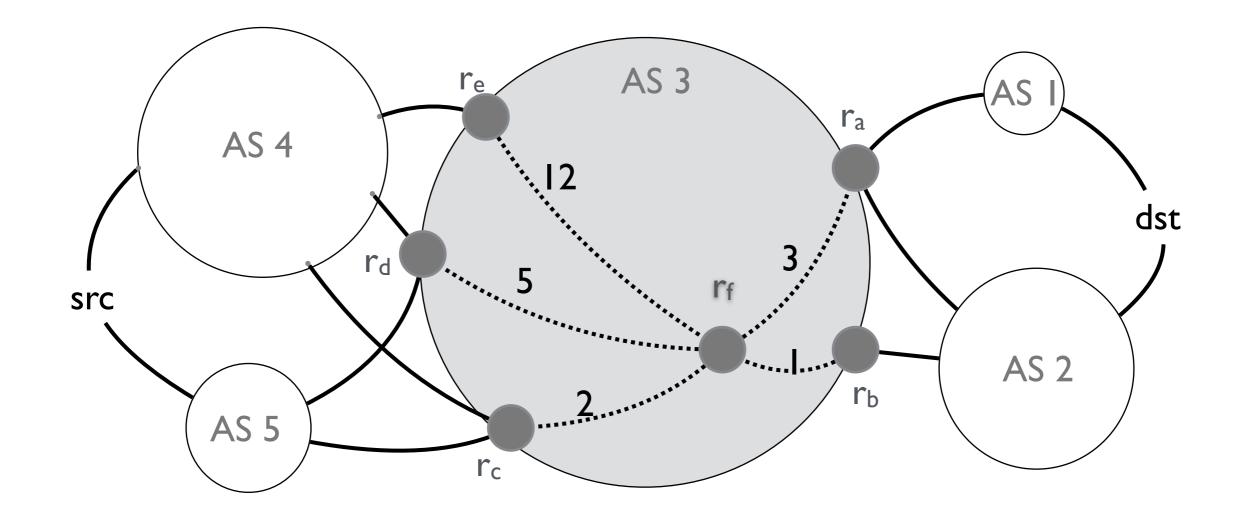
BGP and AS policy

(example) hot potato policy of AS 3 – select a path that minimizes internal cost

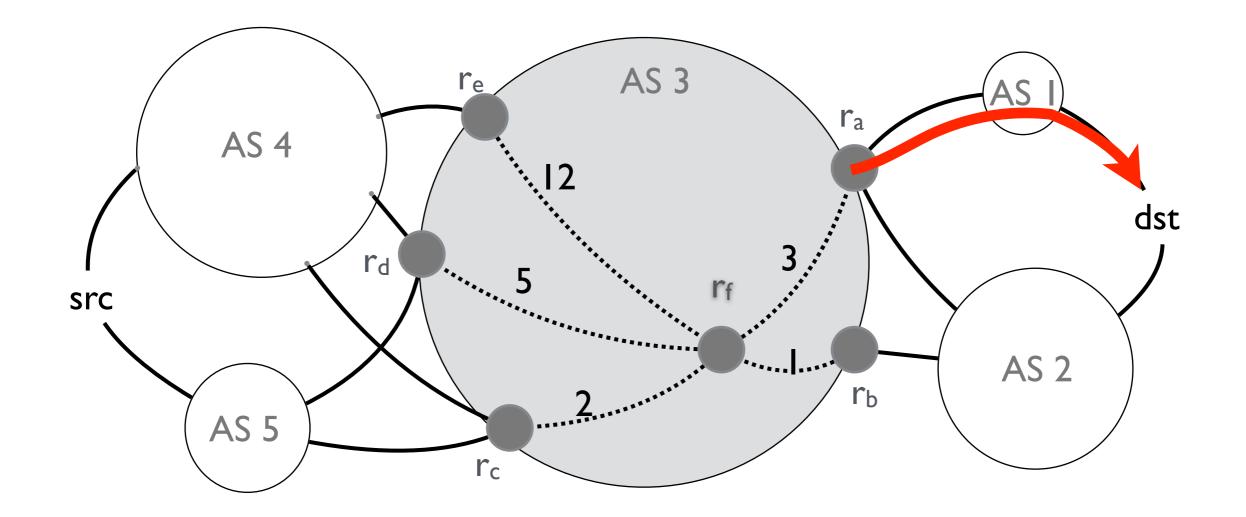


influence policy in the downstream

can AS 4 / AS 5 influence routes in AS3?

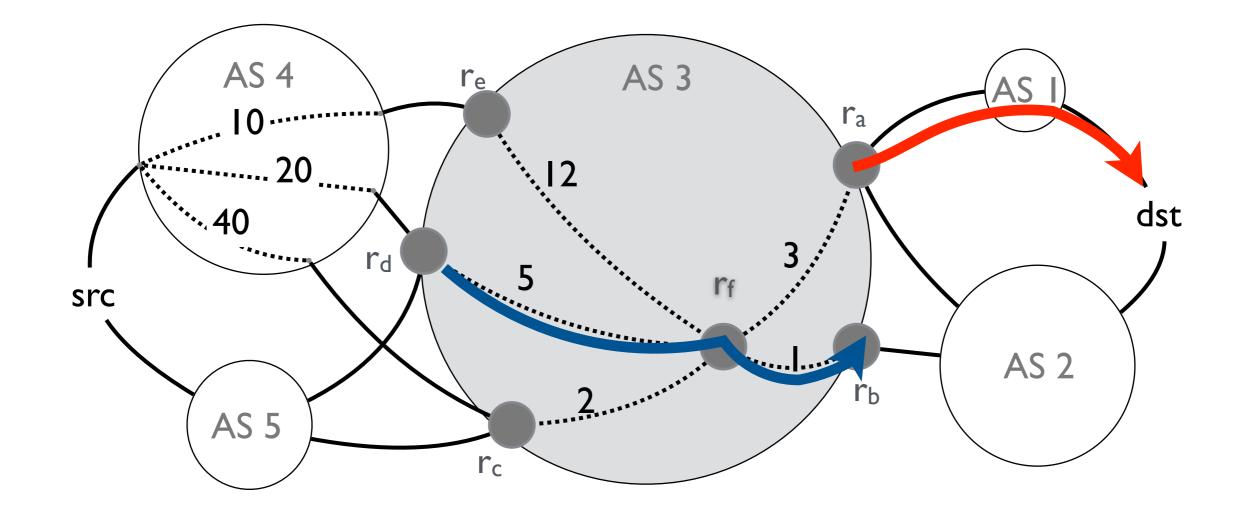


can AS 4 / AS 5 influence routes in AS3?-AS 5 requests AS 3 to bypass AS 2?



can AS 4 / AS 5 influence routes in AS3?

- AS 5 requests AS 3 to bypass AS 2?
- AS 4 demands AS 3 for joint traffic engineering?

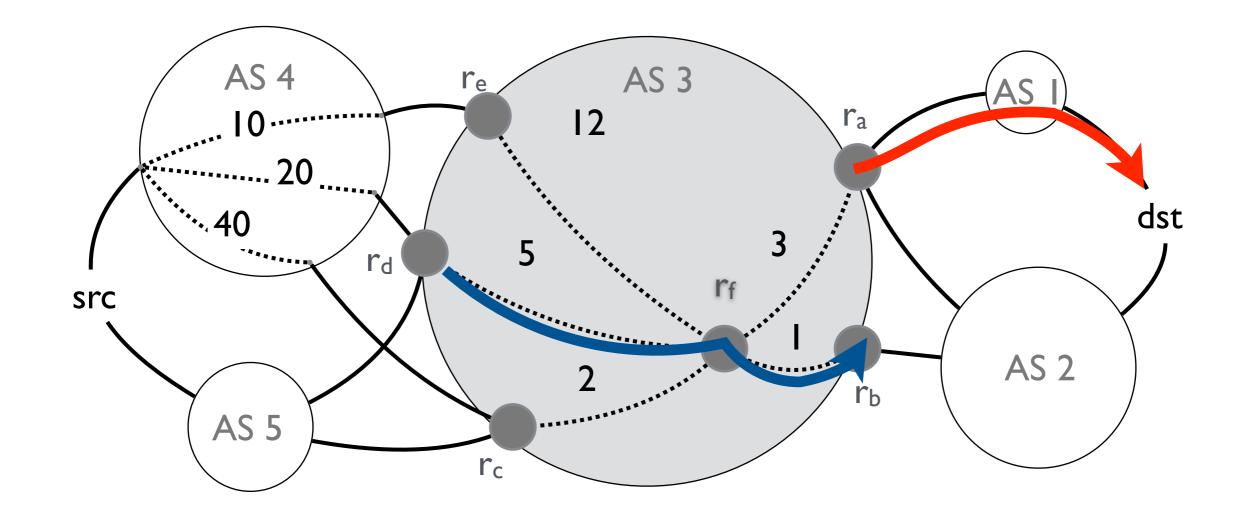


- but the flow of policy attribute in BGP is unidirectional
 - -policy carried by the routes

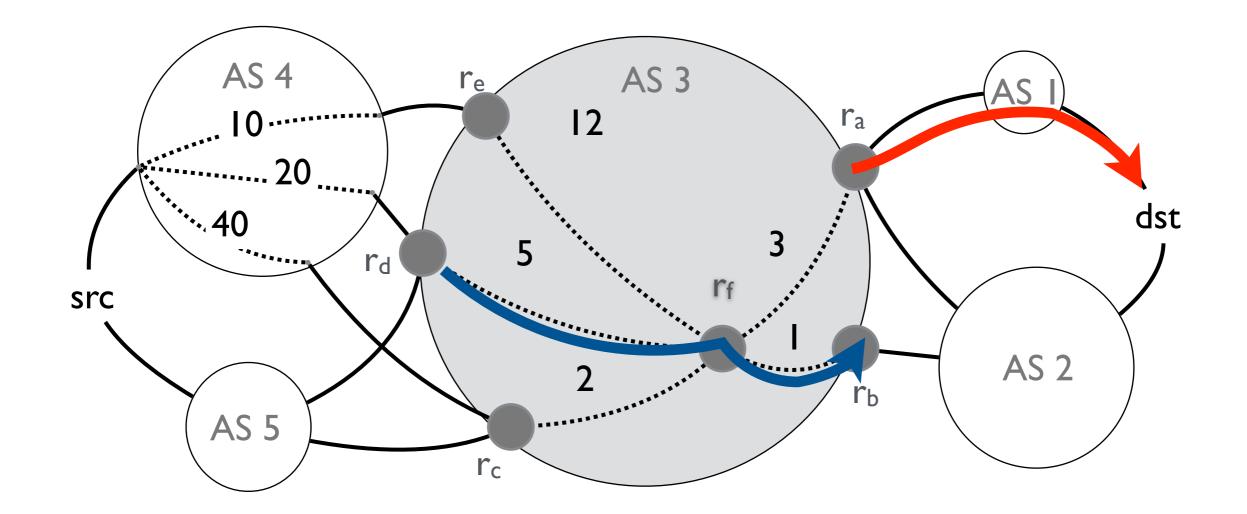
BGP policy is restricted new extensions to path-vector?

can AS 4 / AS 5 influence routes in AS3?

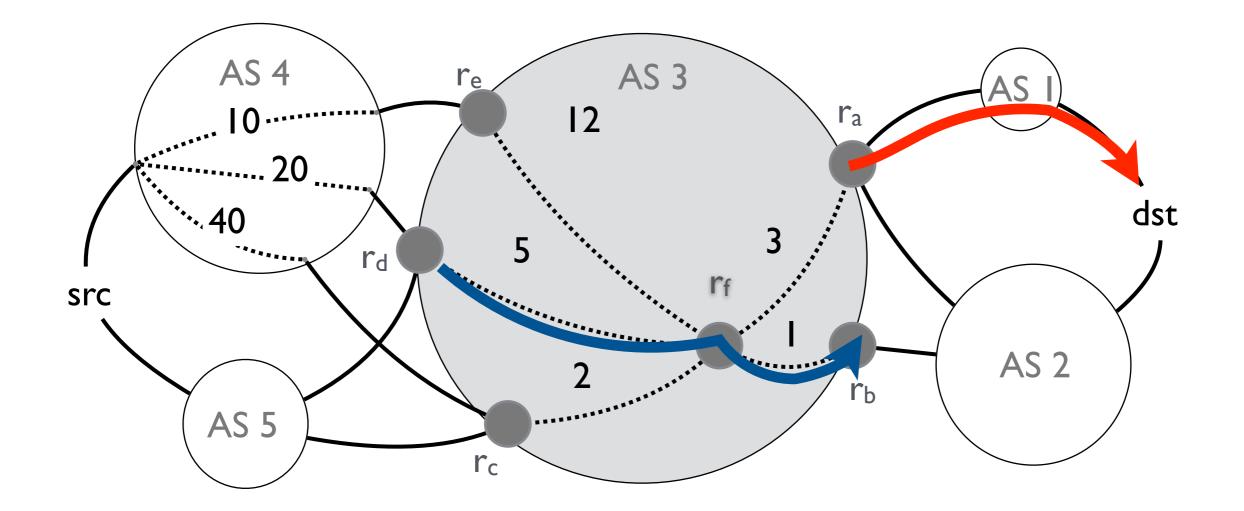
AS 5 ... — BGP + negotiation [MIRO, sigcomm'08]
AS 4 ... — BGP + new attribute [Wiser, sigcomm'07]



simultaneously, AS 5 wants to avoid AS 2, AS 4 demands joint TE?

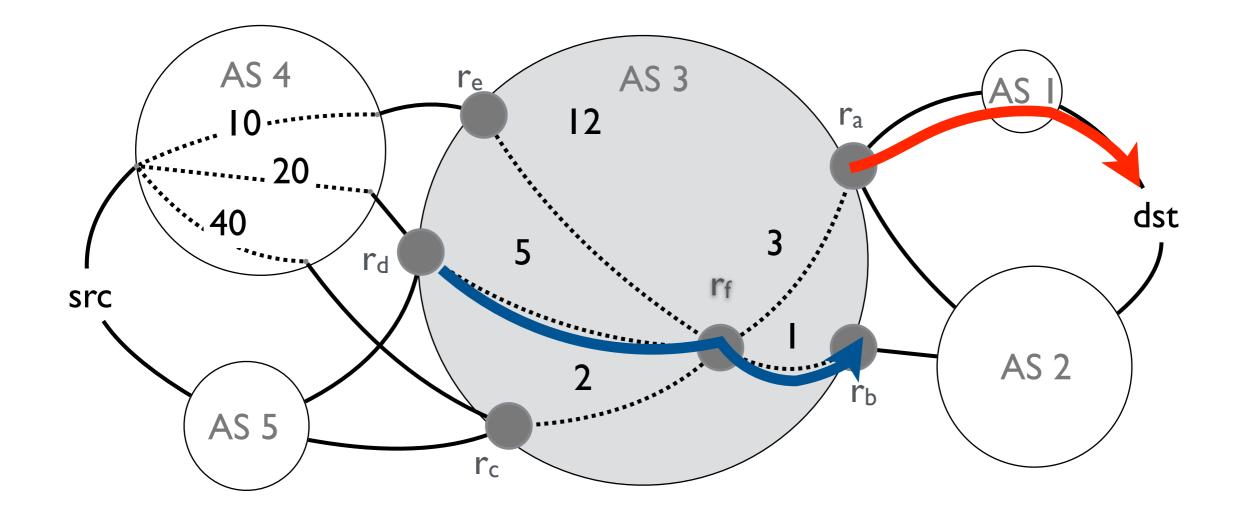


still impossible with (BGP + MIRO + Wiser)

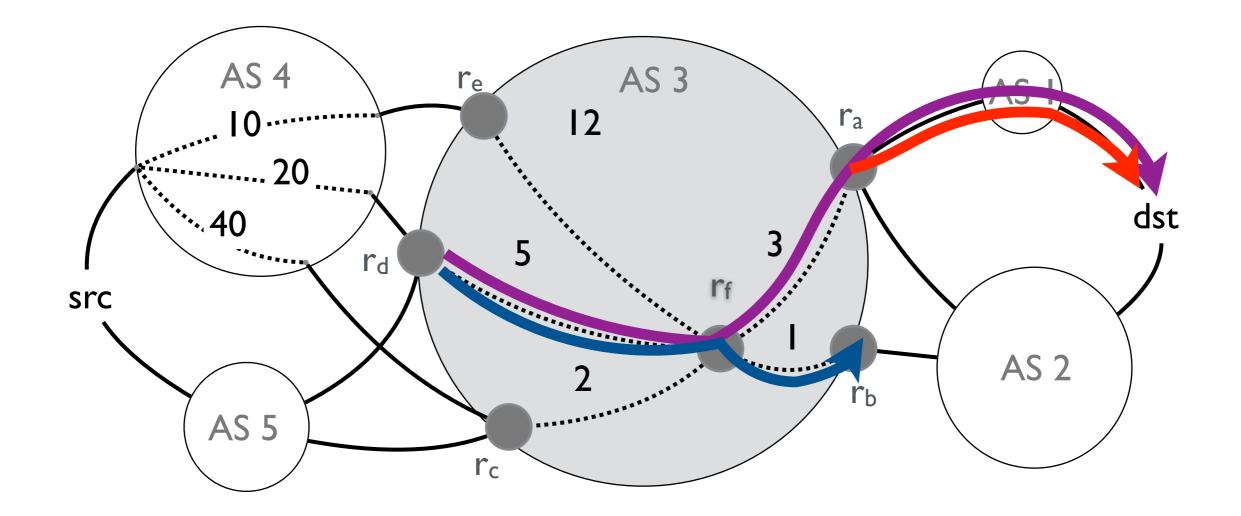


(AS 5 avoids AS 2)+(AS 4 demands joint TE)

- AS 3 needs to *properly* combine the sub-routes: simple concatenation does not work!

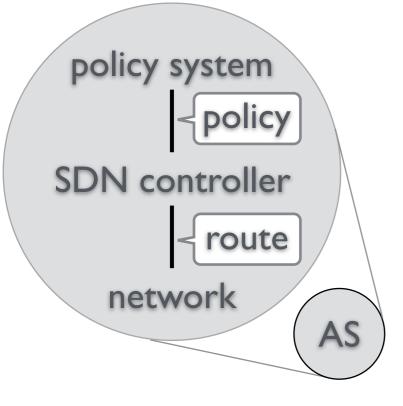


- (AS 5 avoids AS 2)+(AS 4 demands joint TE)
 - -AS 3 needs to resolve conflicts modify the subpath by the less important policy (AS 4 demands joint TE)



coupling routes and policies — including any path vector based policy (BGP, MIRO, Wiser ...) — is inherently flawed

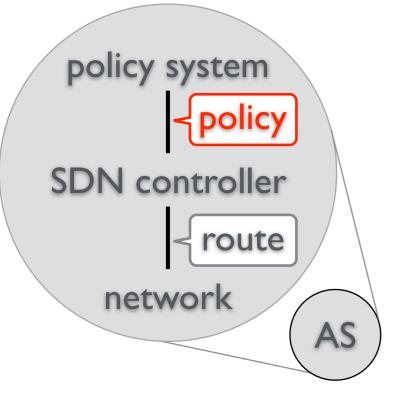
separate policies from routes



a new policy system with SDN

- SDN controller: route discovery and dissemination
- policy system: express and process high-level intention

separate policies from routes



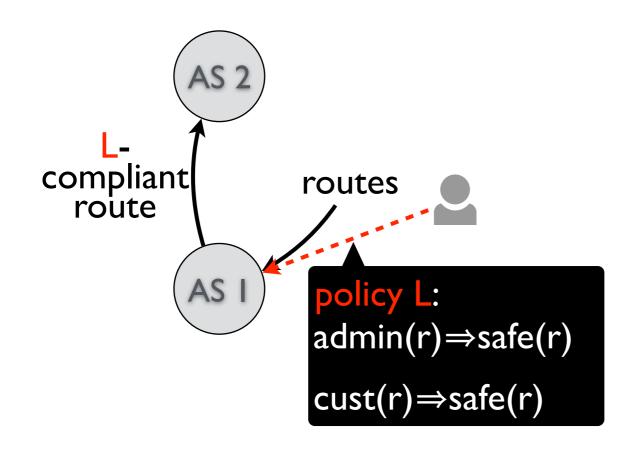
a new policy system with SDN

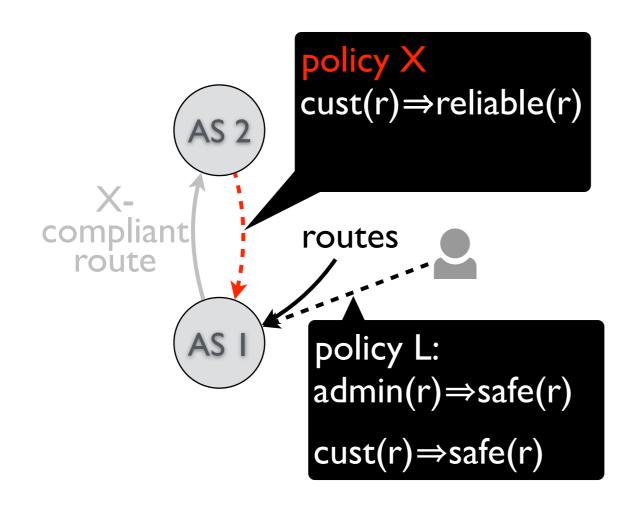
- SDN controller: route discovery and dissemination
- policy system: express and process high-level intention

key idea: making policies logic statements that like routes — freely flow and interact

more flexible

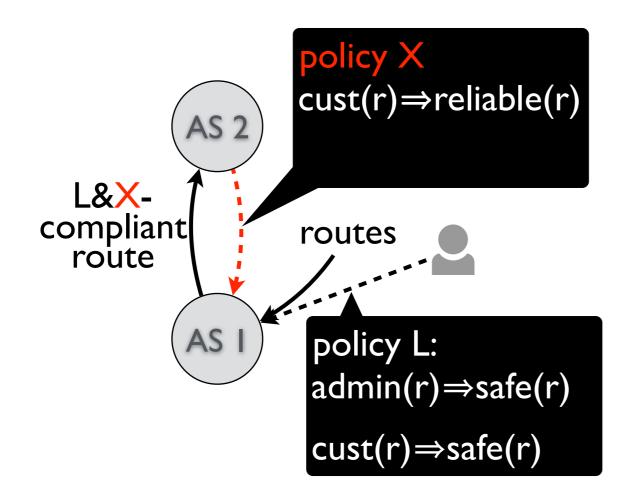
 logic unifies disparate policies





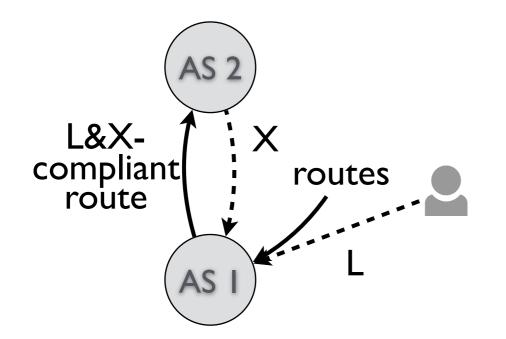
more flexible

- logic unifies disparate policies
- immediately allows
 control of the
 downstream



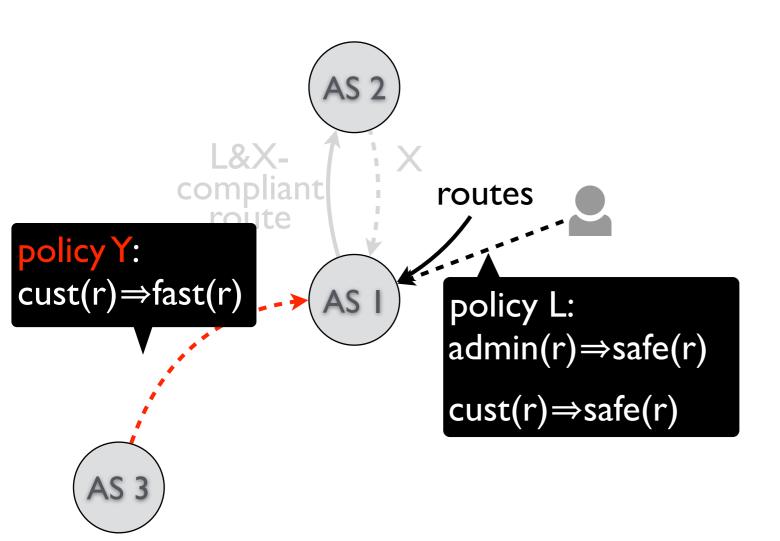
more flexible

- logic unifies disparate policies
- immediately allows
 control of the
 downstream



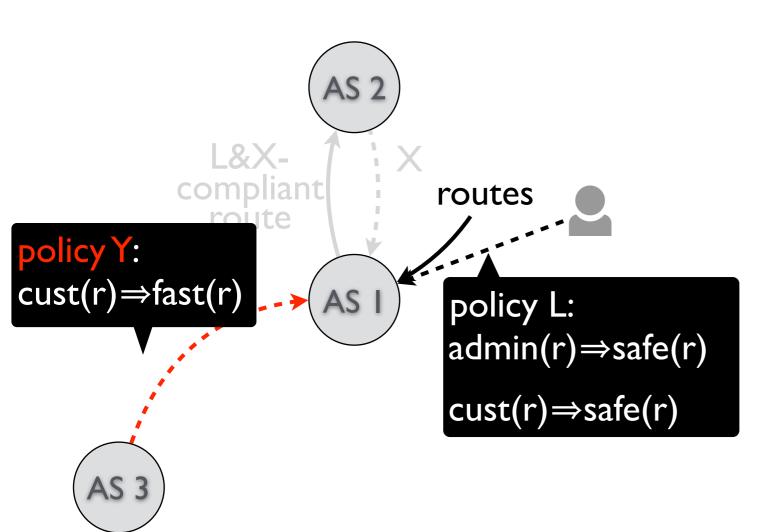
more flexible

- logic unifies disparate policies
- immediately allows
 control of the
 downstream



coordination

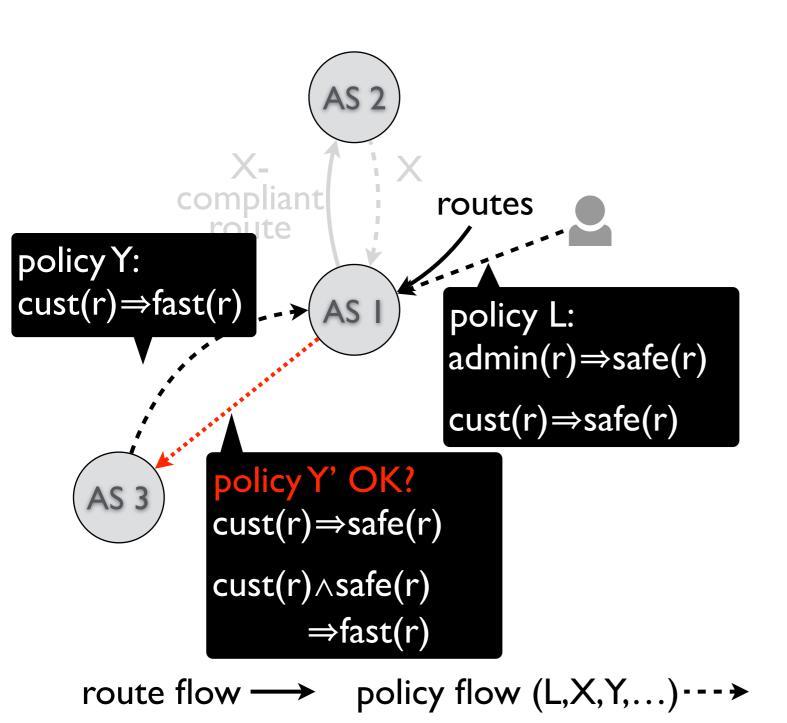
 compute the "impact" of a policy on another policy



coordination

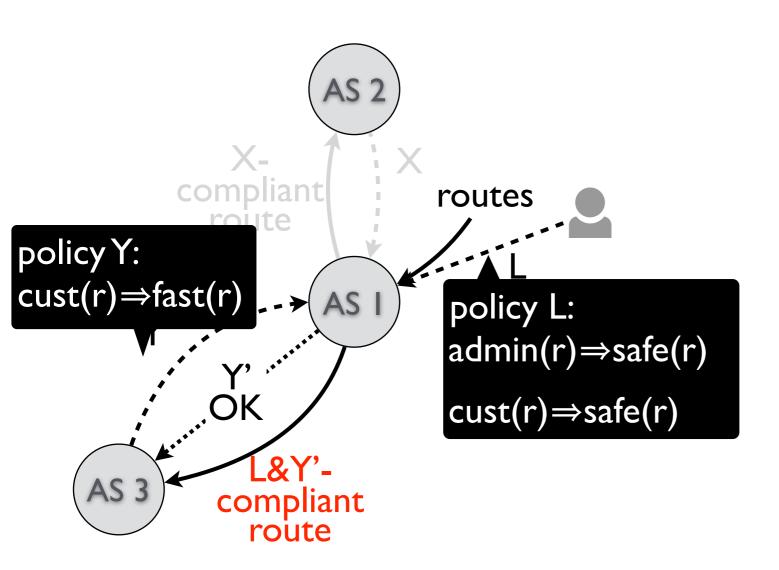
 compute the "impact" of a policy on another policy

but some routes cannot be simultaneously safe and fast route flow → policy flow (L,X,Y,...)····>



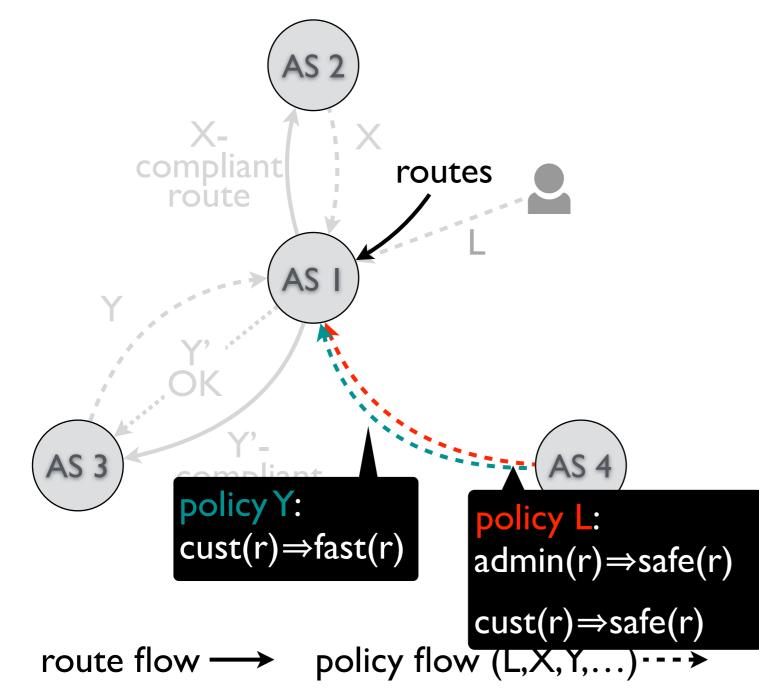
coordination

 compute the "impact" of a policy on another policy



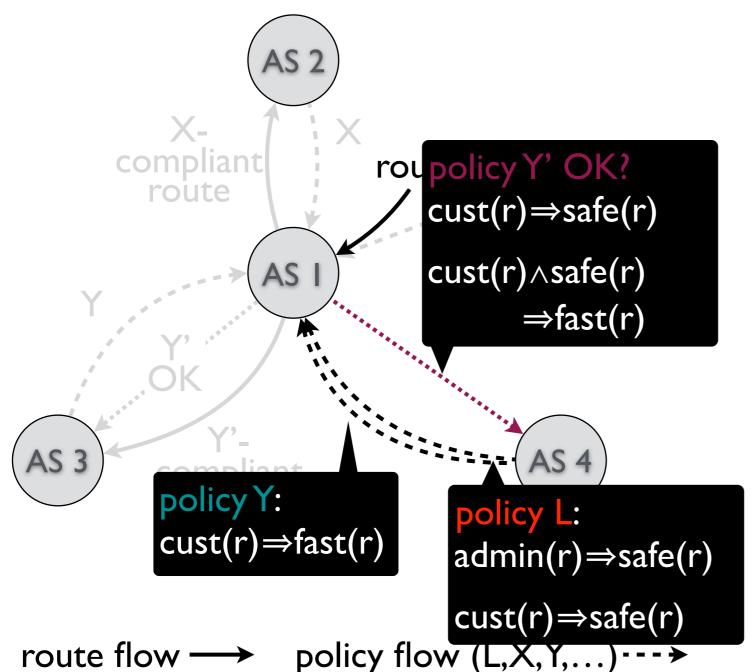
coordination

 compute the "impact" of a policy on another policy



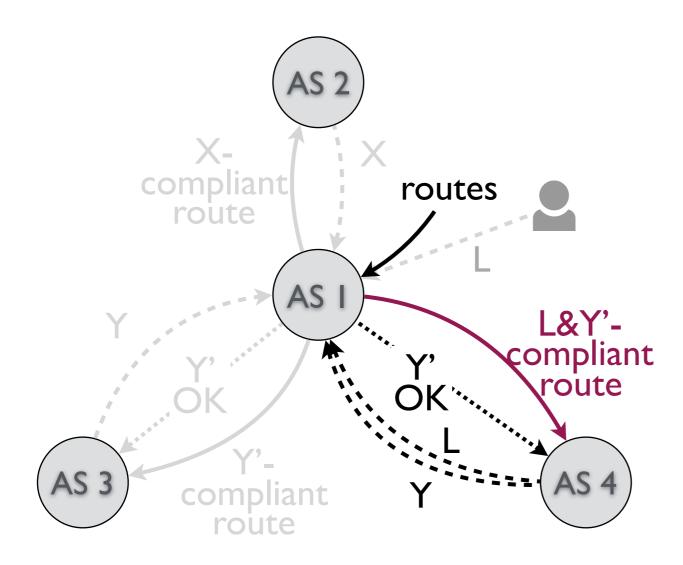
coordination

 compute the "impact" of a policy on another policy



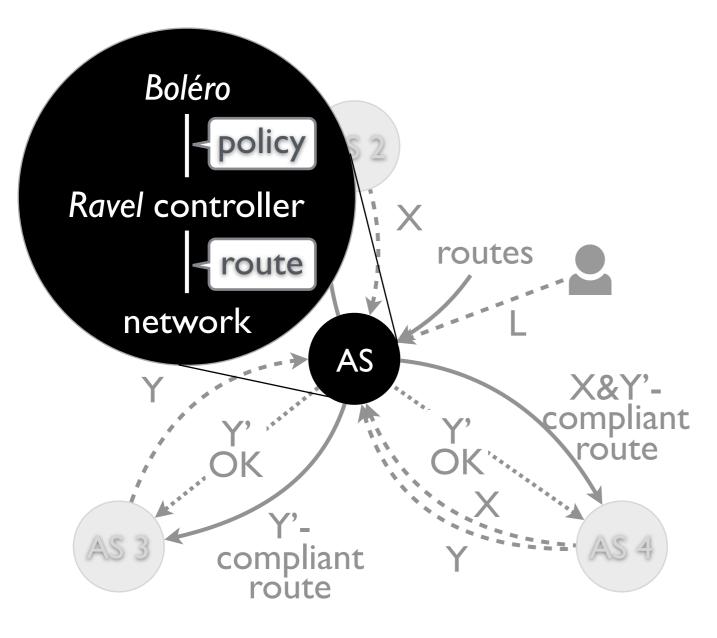
coordination

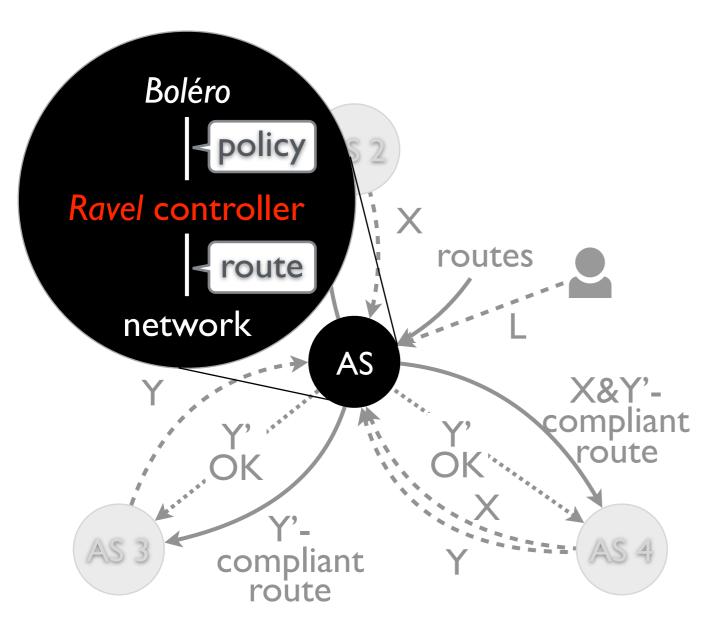
 compute the "impact" of a policy on another policy



coordination

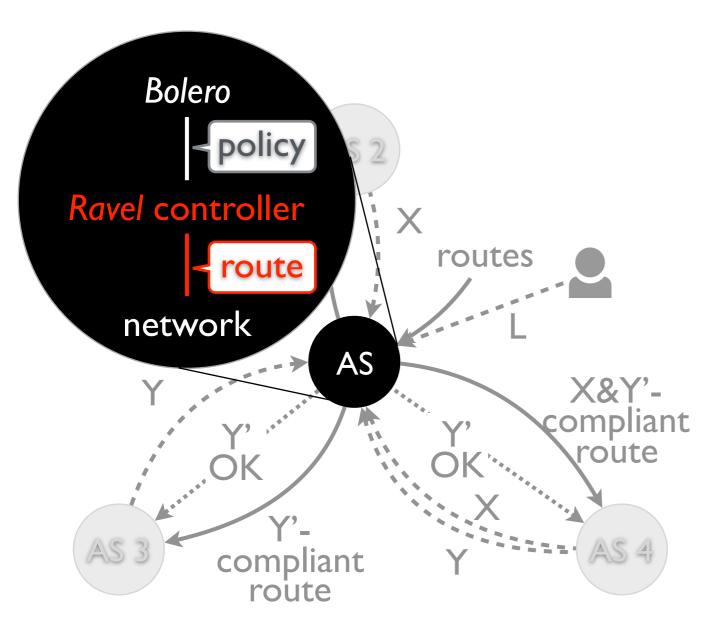
 compute the "impact" of a policy on another policy





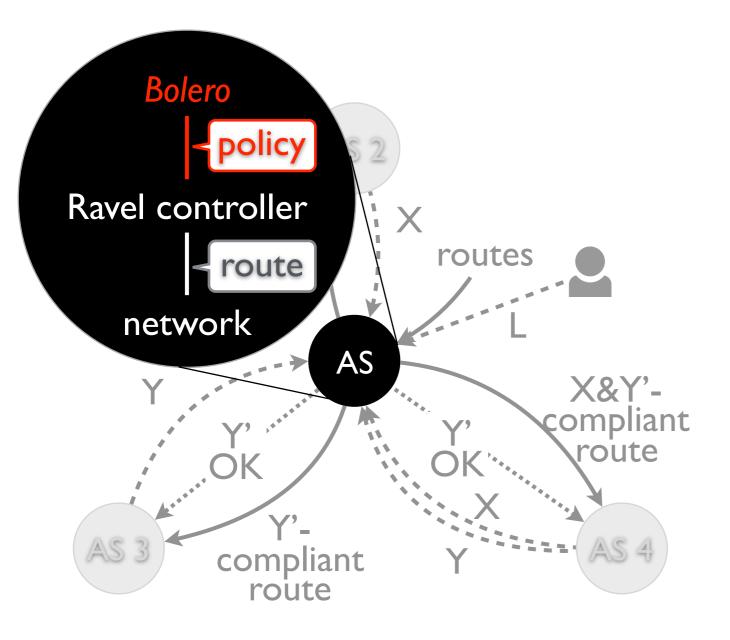
Ravel

 uses a high-performance database as the controller



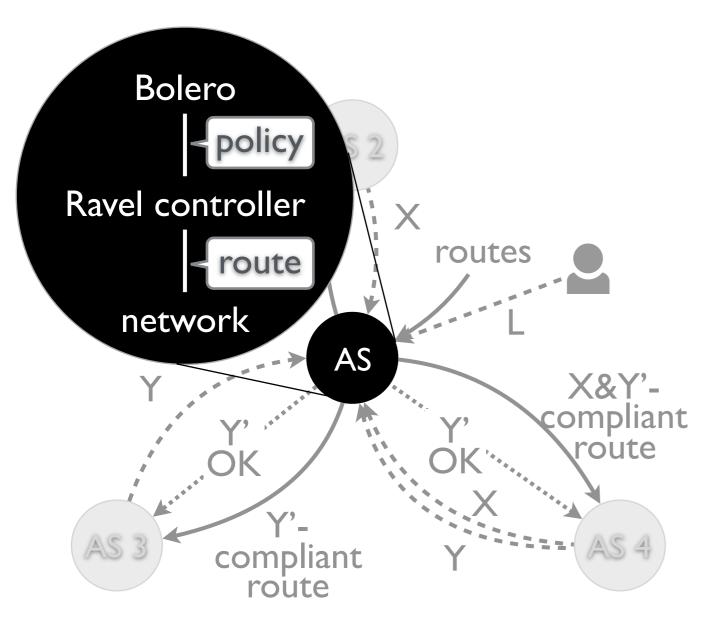
Ravel

- uses a high-performance database as the controller
- discovers and disseminates routes (maintain database tables)



Ravel controller

- uses a high-performance database as the controller
- discovers and disseminates BGP routes
- Boléro policy system
 - exchanges and processes
 logic policies

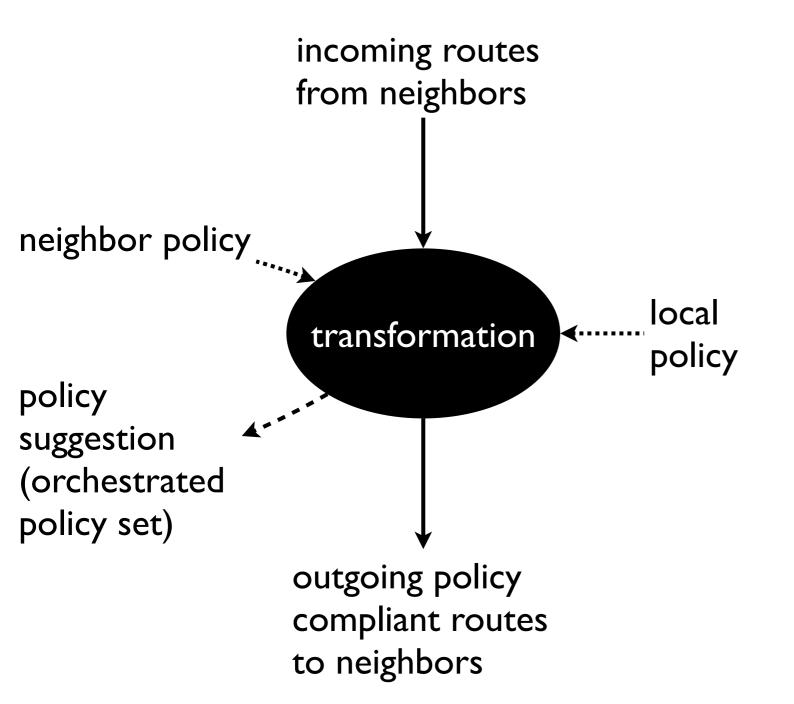


Ravel controller

- uses a high-performance database as the controller
- discovers and
 disseminates BGP routes
- Boléro policy system
 - exchanges and processes
 logic policies

key idea: process policy — semantic knowledge — by semantic transformation

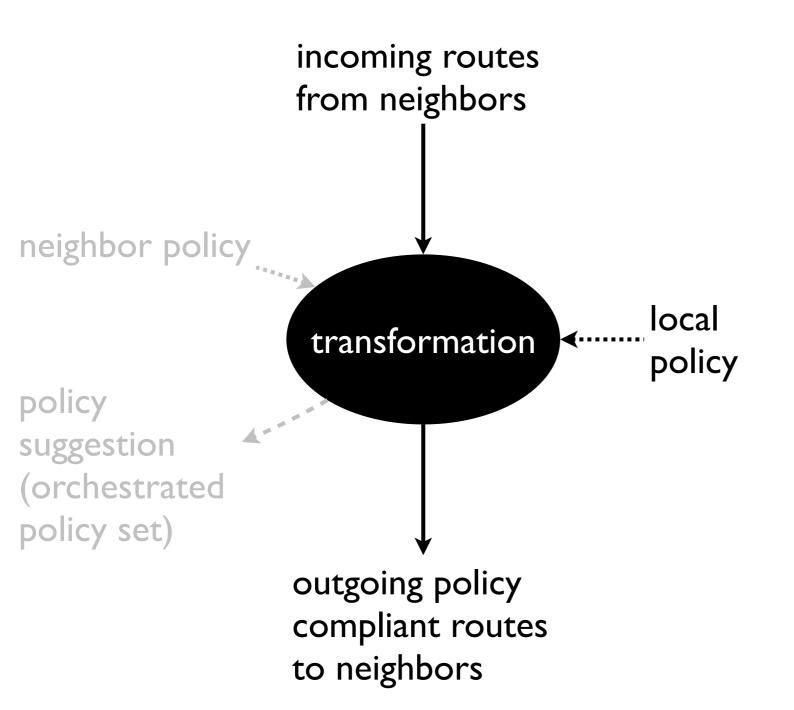
Boléro overview



transformation

- flexible policy routing
- policy coordination

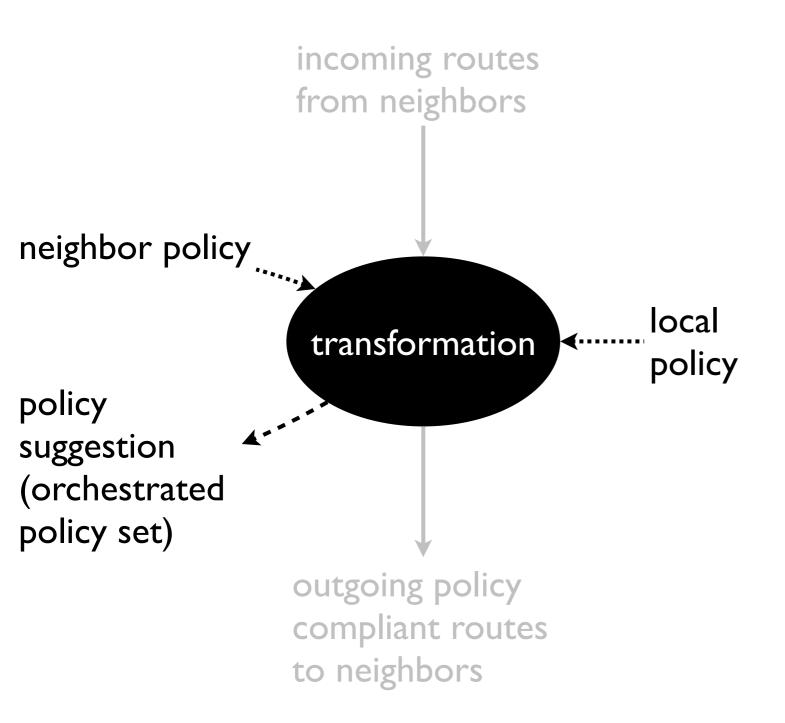
Boléro overview



transformation

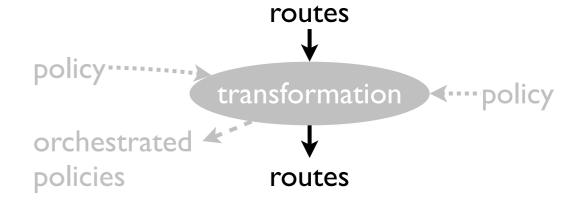
- policy routing
 - transform routes using the policies

Boléro overview



transformation

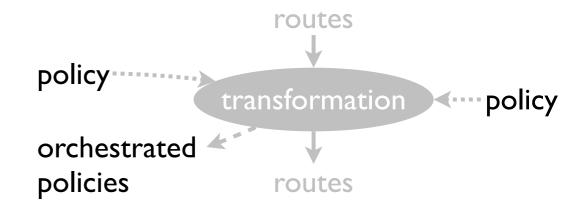
- policy routing
 - transform routes using the policies
- policy coordination
 - transform one by another policy



route

prefix	next hop	AS path	

routing state as queryable tables *factual data*



route

prefix	next hop	AS path	

- --- MIRO-like policy
- :- route(D,N,P), ('AS2' in P).

- routing state as queryable tables
 - factual data
- policy as data integrity constraint (IC)
 - logical statement about what (must be avoided) are the valid route data
 - semantic data

policy transformation policy orchestrated policies routes

route

prefix	next hop	AS path	

- --- MIRO-like policy
- :- route(D,N,P), ('AS2' in P).

meaning

False <--

route(I,R,P)/ $\{ ('AS2' in P) \}$

- routing state as queryable tables
 - factual data
- policy as data integrity constraint (IC)
 - logical statement about what (must be avoided) are the valid route data
 - semantic data

policy transformation policy orchestrated policies routes

route

prefix	next hop	AS path	

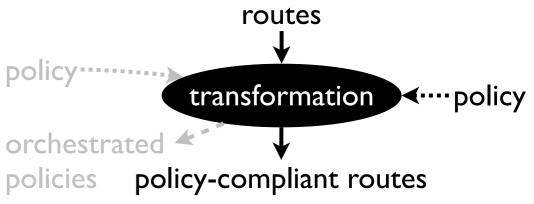
- --- MIRO-like policy
- :- route(D,N,P), ('AS2' in P).
- --- Wiser-like policy
- :- route(D,N,P),

```
Wiser(D,R,C),Advertise(R,C<sub>2</sub>),
Wiser(D,R',C'),Advertise(R',C<sub>2</sub>'),
```

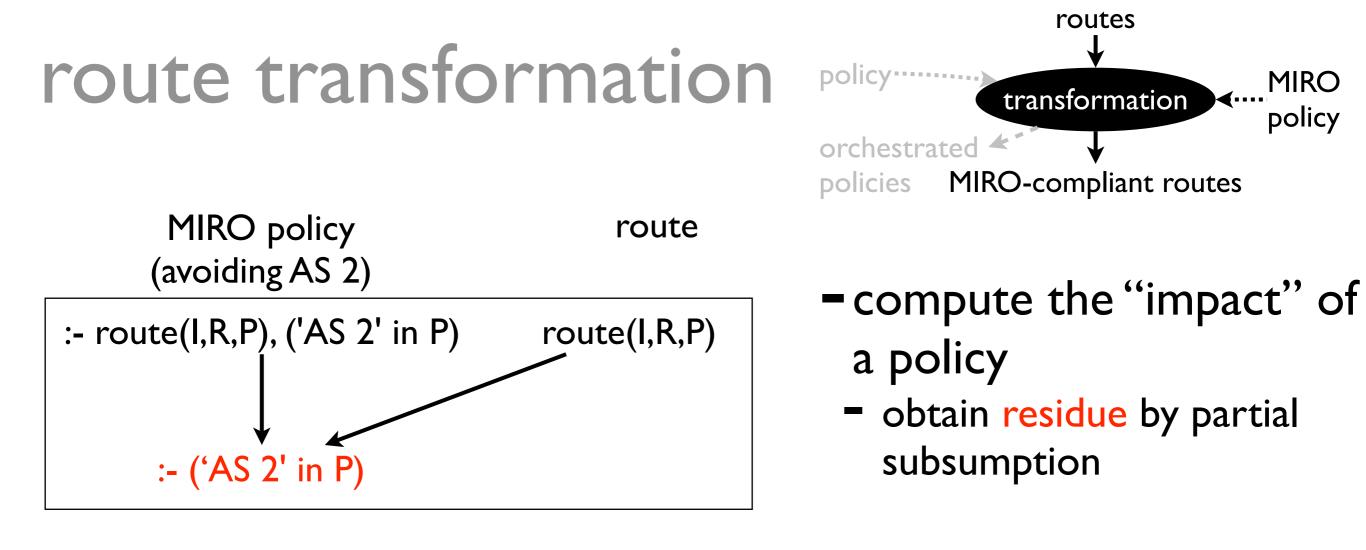
```
C+C2>C'+C2'.
```

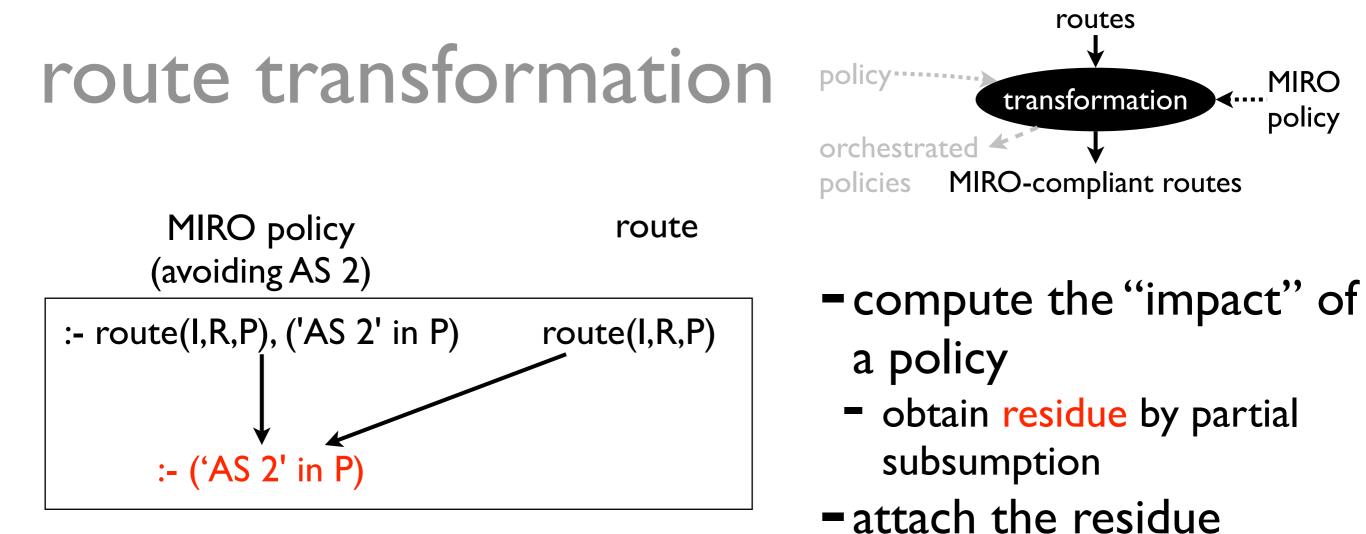
- routing state as queryable tables
 - factual data
- policy as data integrity constraint (IC)
 - logical statement about what (must be avoided) are
 the valid route data semantic data

route transformation policy.



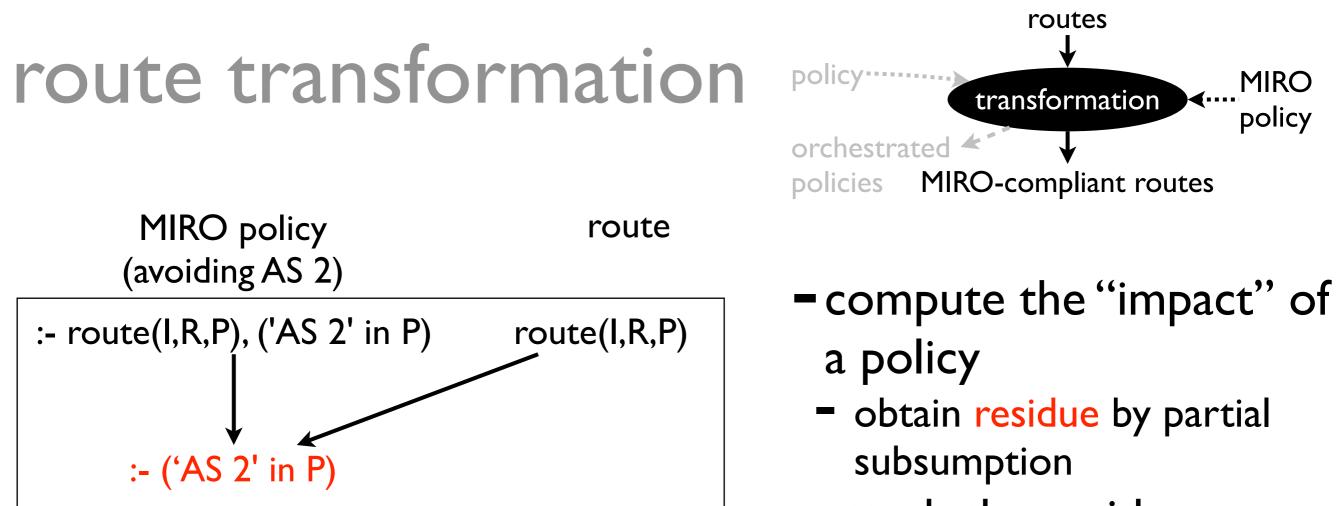
- compute the "impact" of a policy





MIRO-compliant route

route(I,R,P) {:-('AS2' in P)}



-attach the residue

MIRO-compliant route

route(I,R,P) {:-('AS2' in P)}

meaning
{¬('AS2' in P)}

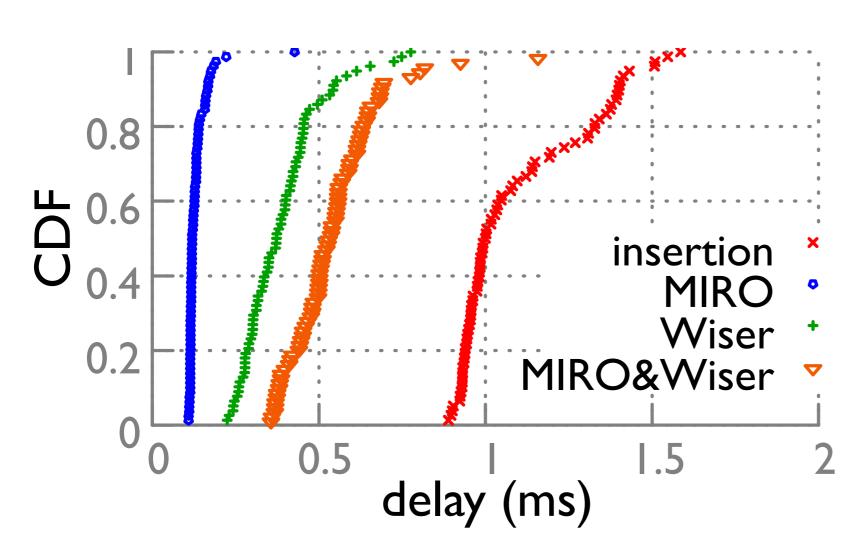
routes policy transformation Wiser transformation MIRO MIRO-compliant routes Wiser policie Wiser policy - compute the "impact" of :- route(D,N,P), a policy Wiser(D,R,C), Advertise(R,C₂), - attach the residue Wiser(D,R',C'), Advertise(R',C₂'), C+C2>C'+C2'.

- --- MIRO-compliant Wiser policy
- :- route(D,N,P),

```
Wiser(D,R,C),Advertise(R,C<sub>2</sub>),
Wiser(D,R',C'),Advertise(R',C<sub>2</sub>'),
C+C2>C'+C2', {¬('AS2' in P)}.
```

preliminary result

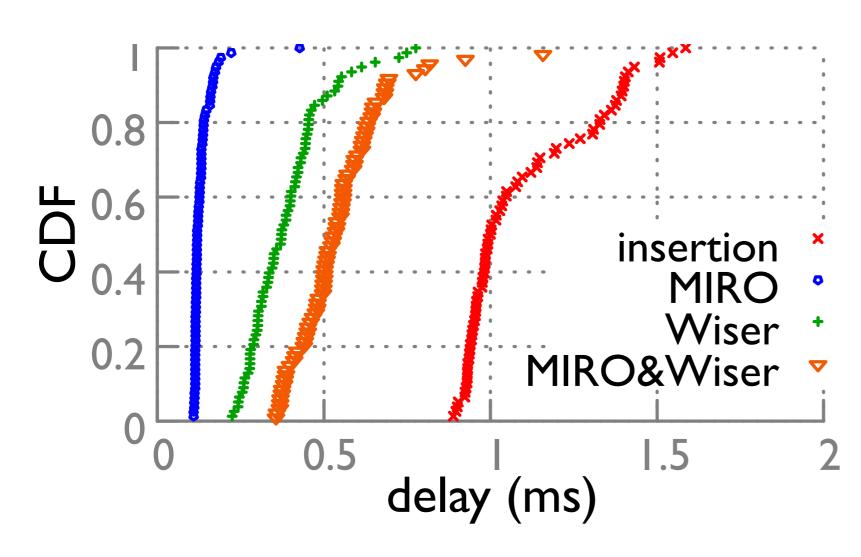
time the database delay for computing policycompliant routes



- policies
 - MIRO
 - Wiser
 - MIRO&Wiser
- topology: Pocketful
 - ISP topology
 - embedded in Skitter
 - AS-level topology
- incoming routes
 - Routeview BGP feeds

preliminary result

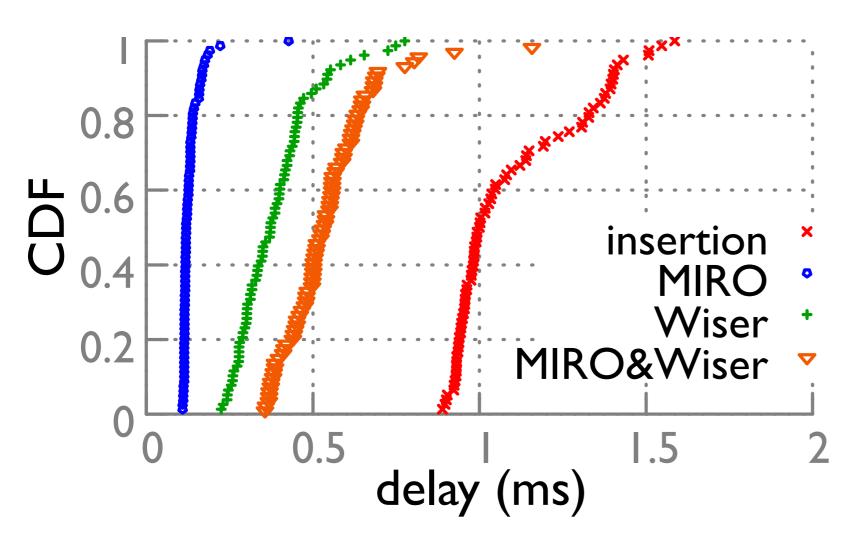
time the database delay for 10,000 BGP feeds



- delay is small and scales well
 - 95% route insertion <
 1.424ms
 - 95% MIRO < .174ms
 - **-** 95% Wiser < .64ms
 - 95% Wiser&MIRO < .
 844ms

preliminary result

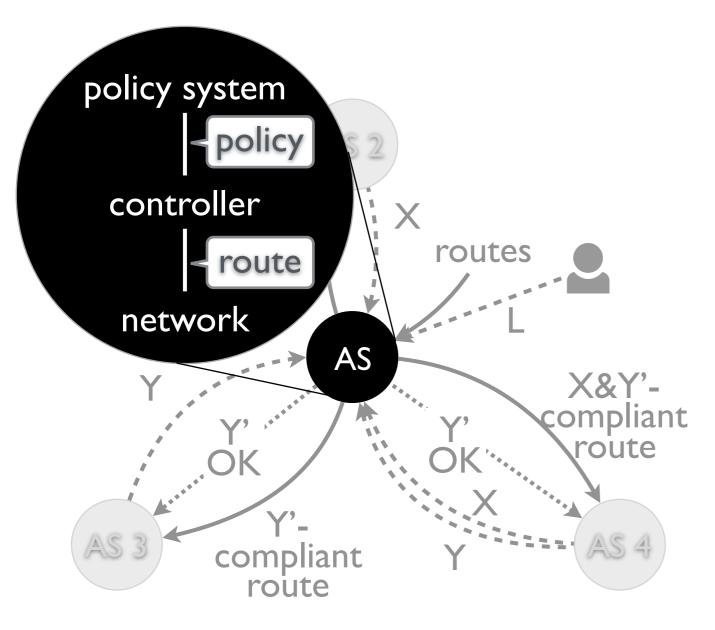
time the database delay for 10,000 BGP feeds



- delay is small and scales well
- delay grows as policy becomes more complex

conclusion

coupling routes and policies is inherently flawed



this talk

decouple policies from routes
 by a new policy system with
 SDN

benefits

- flexible policies
- automated coordination

future work

- implementation
- anonymize policies