

A Throughput-Centric View of the Performance of Datacenter Topologies

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When experts design a network, they try to provision the network to handle expected traffic demands...

When cloud providers design a datacenter network, they try to provision the network to handle any possible traffic demand.

* To a first approximation. We discuss oversubscription in the paper.

Why any possible traffic demand

Datacenters are long-lived

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Traffic can change significantly

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Any feasible traffic demand

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Datacenters are long-lived

Traffic can change significantly

Any feasible traffic demand

Cloud application performance independent of VM placement

Why any possible traffic demand

Data

Traffic

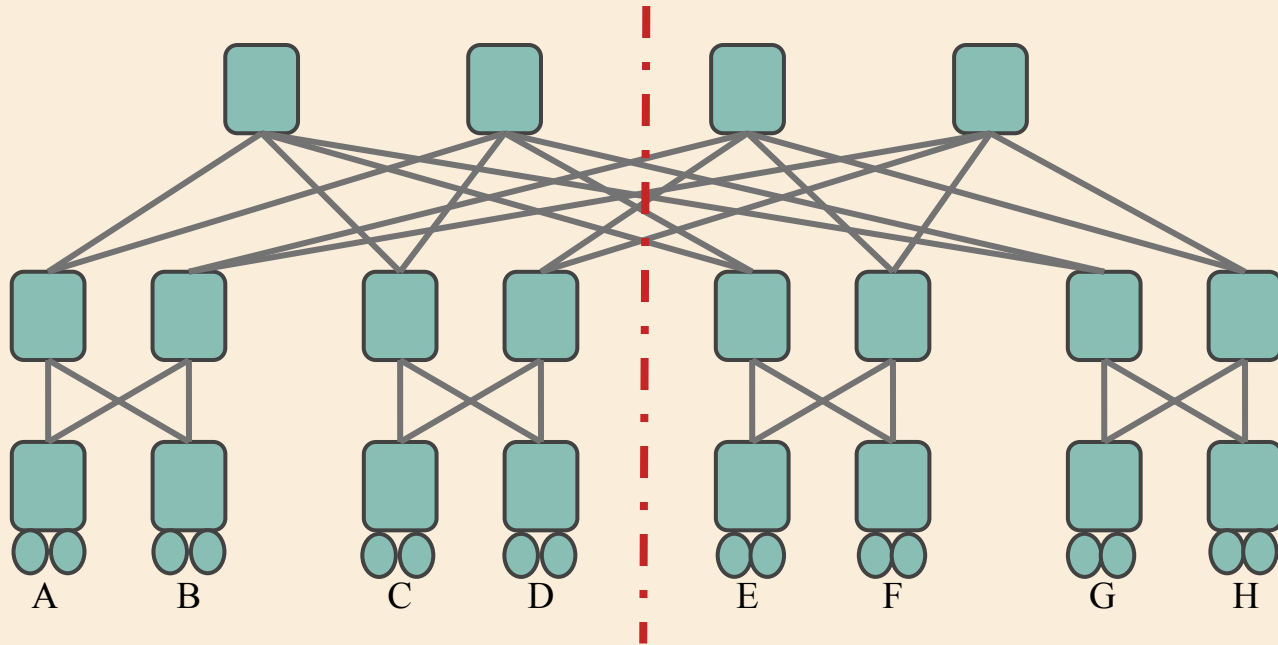
traffic

Non-blocking Topology;
A topology that does not block
any traffic demand

Cloud application performance independent of VM placement

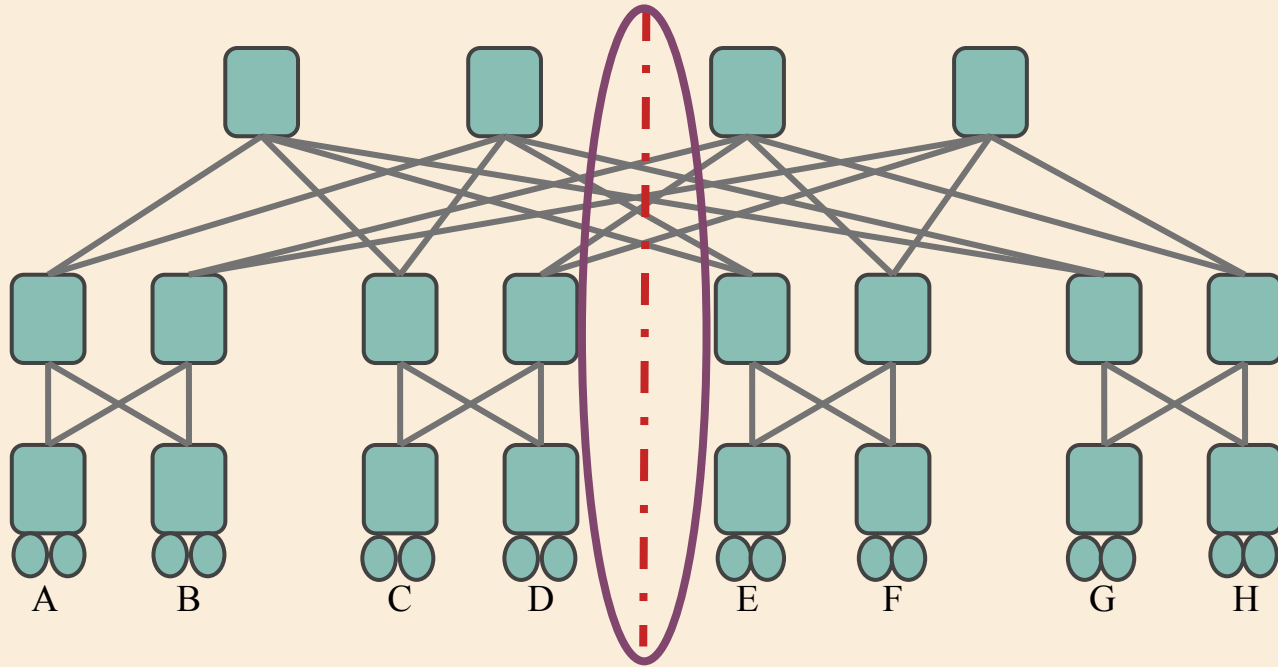
How to assess whether a
datacenter topology is
non-blocking?

Early Work uses Bisection Bandwidth



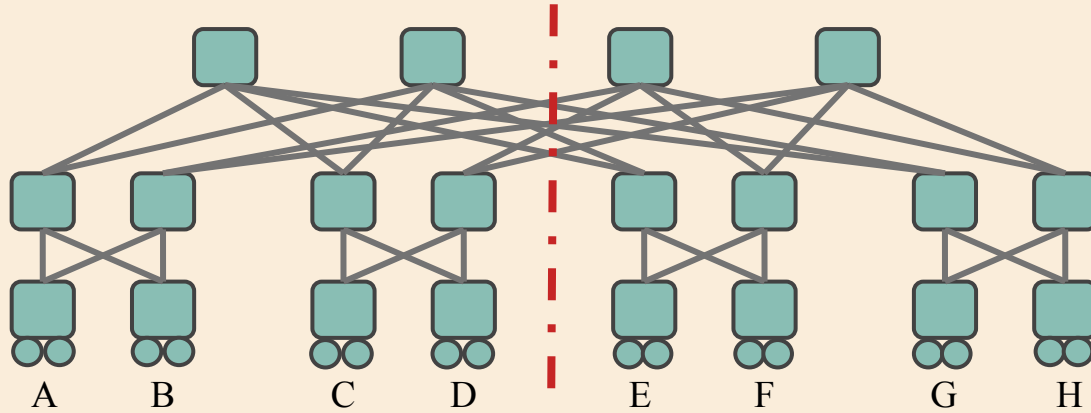
Bisection Bandwidth

Early Work uses Bisection Bandwidth



Bisection Bandwidth

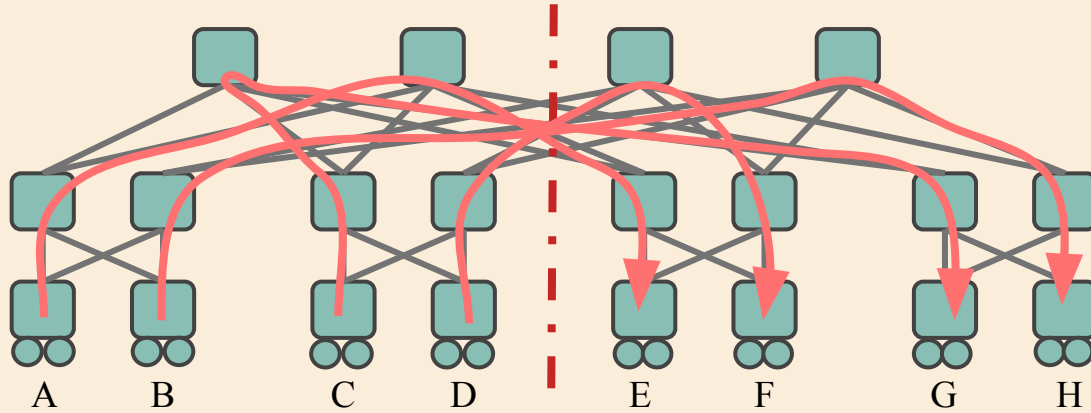
Early Work uses Bisection Bandwidth



Full Bisection Bandwidth

Bisection Bandwidth $\geq \#servers/2$

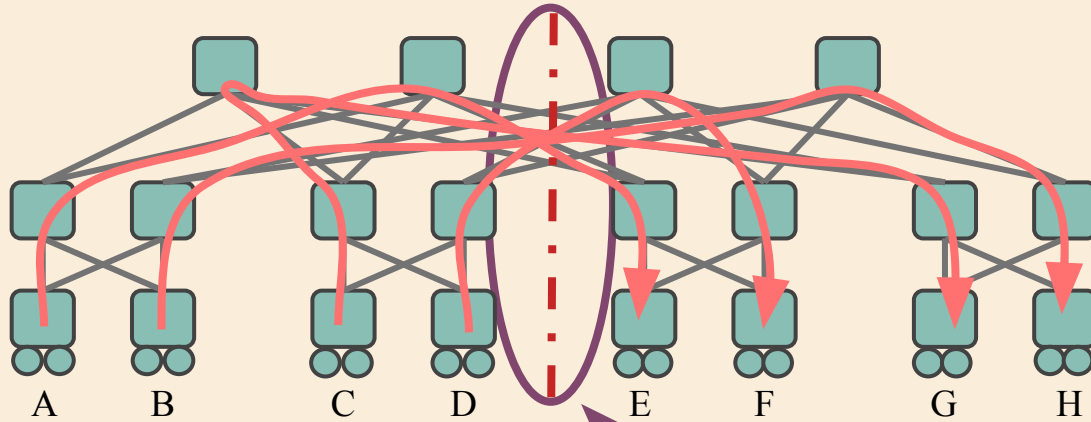
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Full Bisection
Bandwidth



Non-blocking
Topology

Early Work uses Bisection Bandwidth

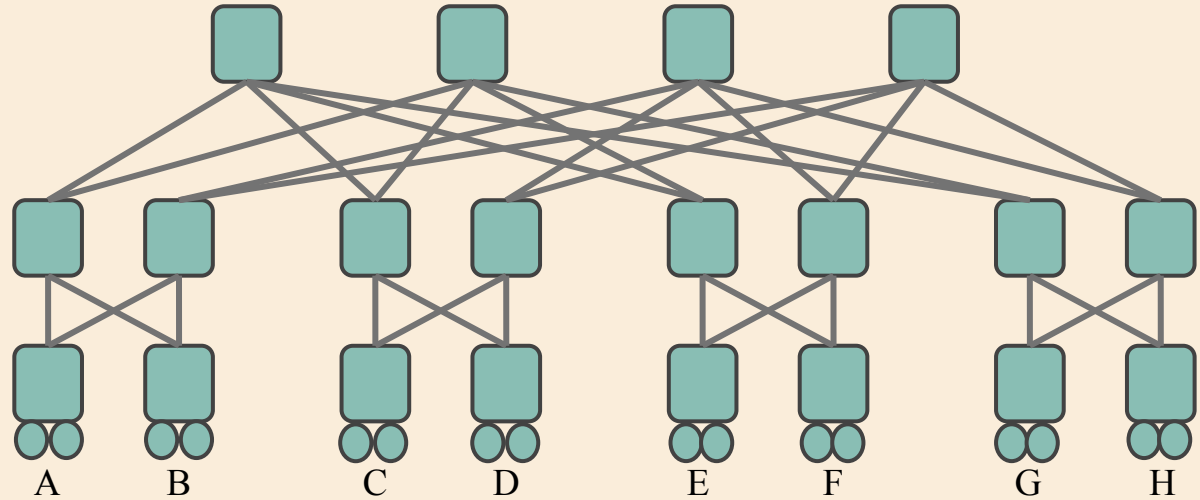
Full Bisection
Bandwidth



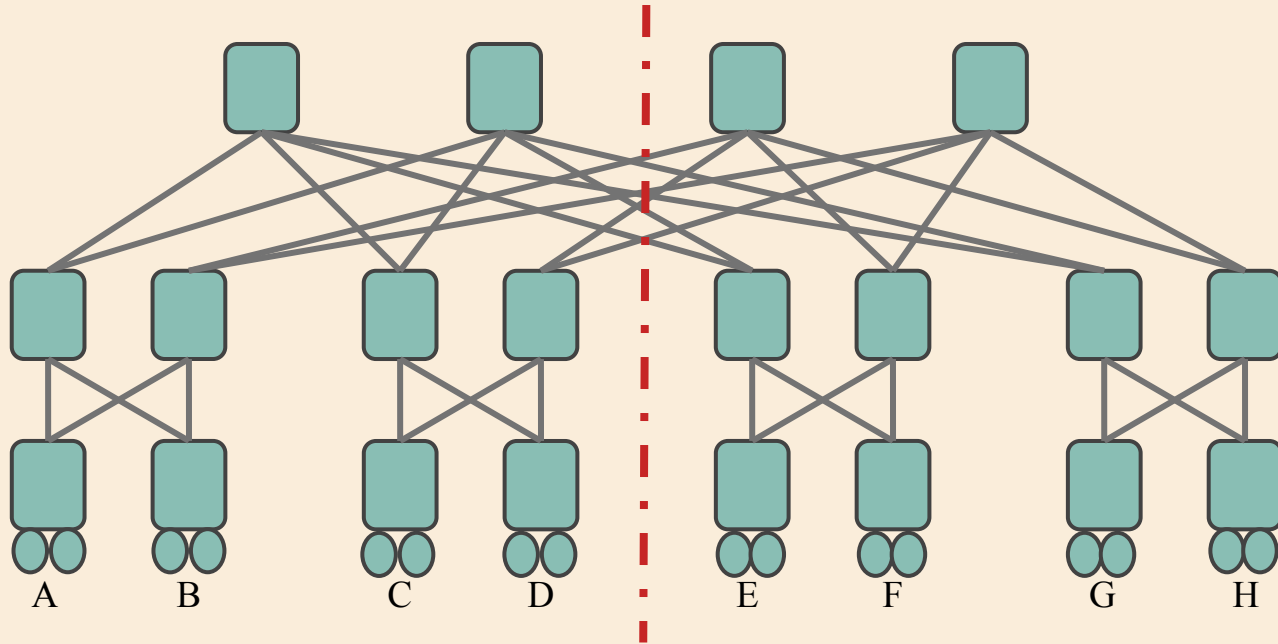
Non-blocking
Topology

This holds for a specific topology family called **Clos**.

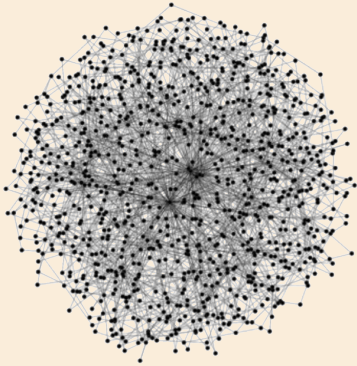
Most Commercial Datacenters are Clos



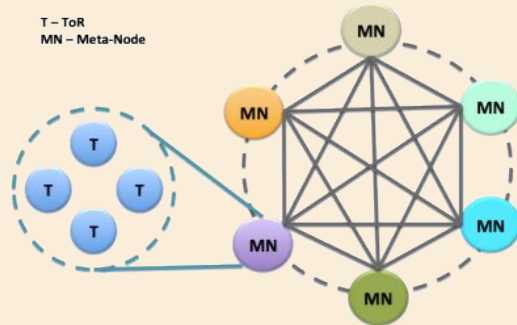
But Clos is Expensive



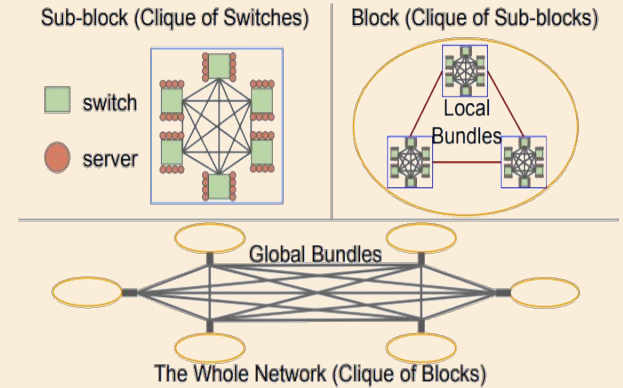
Recently Proposed Topologies: Expanders



Jellyfish
[NSDI'12]



Xpander
[CoNEXT'16]



FatClique
[NSDI'19]

Recently Proposed Topologies: Expanders

Lower Cost (#Switches, #Links, #Racks,)

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Better Management Complexity (Expansion, Wiring,)

Recently Proposed Topologies: Expanders

Lower Cost (#Switches, #Links, #Racks,)

Better Management Complexity (Expansion, Wiring,)

Better Failure Resiliency (Random Failure,)

For expanders, can **bisection bandwidth** help assess whether topology is non-blocking?

* It is for Clos \rightarrow proof in the paper.

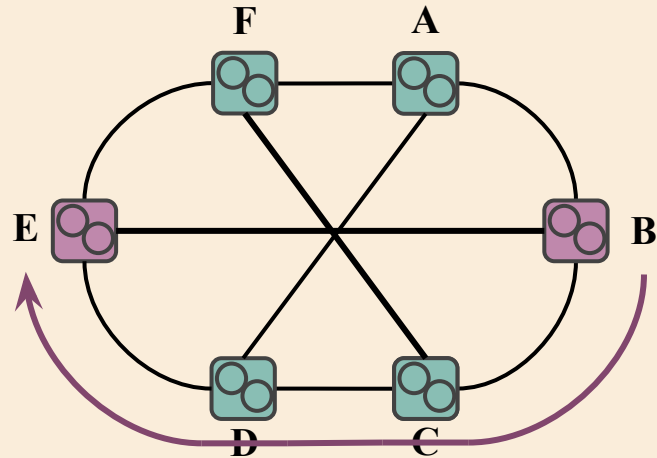
Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain

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Demand from B to E = 2.0

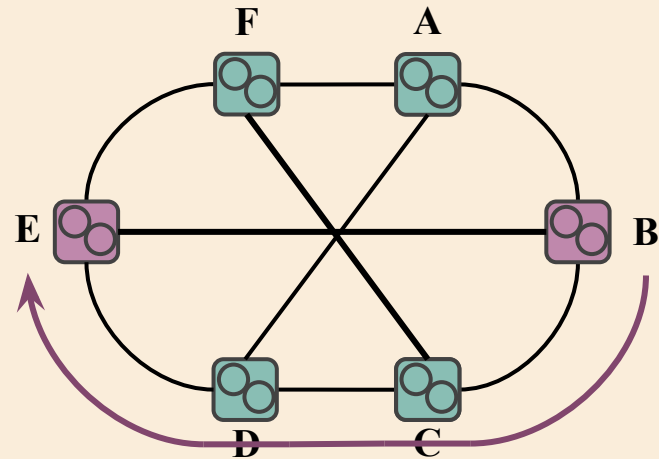


Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain

Demand from B to E = 2.0

Network can sustain = 1.5



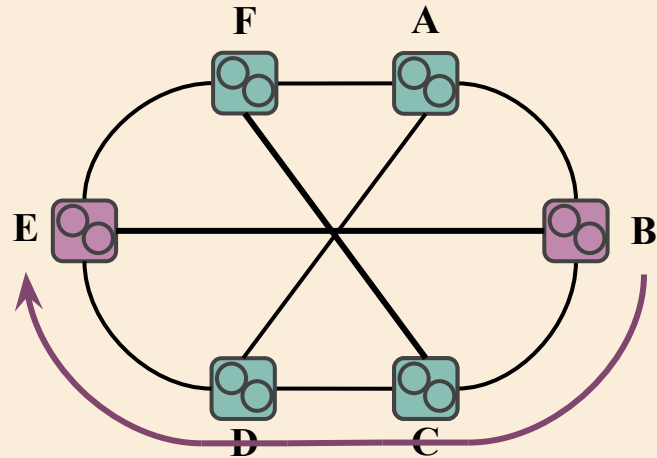
Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain

Demand from B to E = 2.0

Network can sustain = 1.5

Throughput = 0.75



Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain



Throughput of 1 means network can support the traffic matrix

Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain

Throughput of topology is the **smallest throughput** across all possible traffic matrices

Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain

Throughput of topology is the **smallest throughput** across all possible traffic matrices



Throughput of 1 means network is non-blocking

Prior Work Has Proposed Another Metric

Throughput of the topology for a given *traffic matrix* measures the fraction of demand that network can sustain

Throughput of topology is the smallest throughput across all possible traffic matrices

Throughput is expensive to compute

For expanders, is bisection
bandwidth equivalent to
throughput?

Findings

1

A full bisection bandwidth Expander may not have full throughput.

Findings

1

A full bisection bandwidth Expander may not have full throughput.



Theory

There are always exist a size beyond which no full throughput Expander topology exists.

Practice

Even Expanders with 10-15K servers may not have full throughput even if they have full bisection bandwidth

Findings

1

A full bisection bandwidth Expander may not have full throughput.



2

Cost, manageability, and failure resilience comparisons affected significantly when throughput is used at large-scale.

But Computing Throughput is Expensive

An **accurate** upper bound for throughput of Expanders and Clos topologies that **scales** well.

Outline

1

A full bisection bandwidth Expander may not have full throughput.

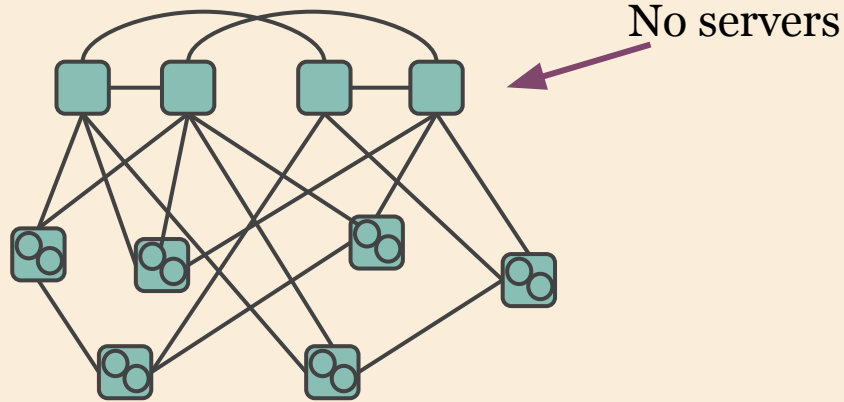
2

Cost, manageability, and failure resilience comparisons affected significantly when throughput is used at large-scale.

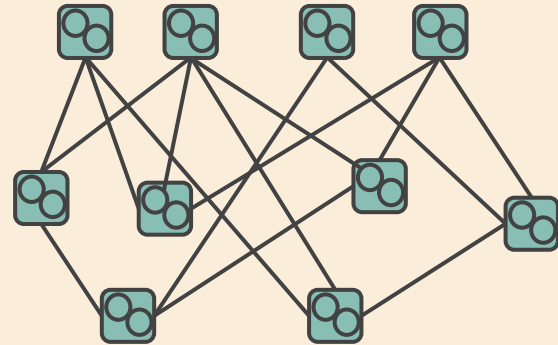
3

An accurate upper bound for throughput of Expanders and Clos topologies that scales well.

Clos vs Expanders



Clos



Expanders

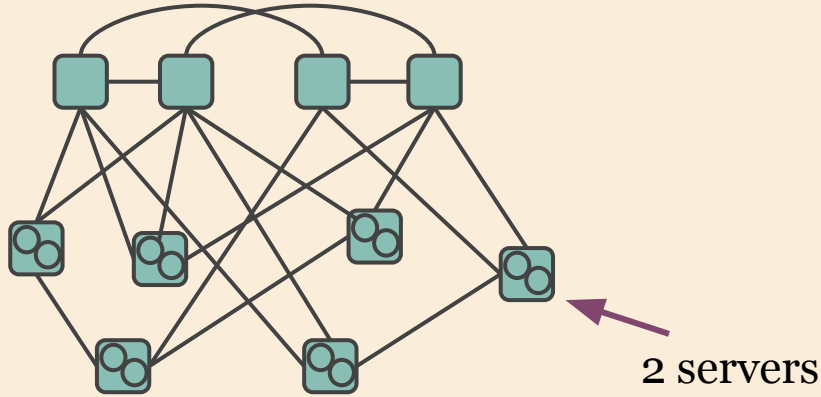


Switch
with 2
servers

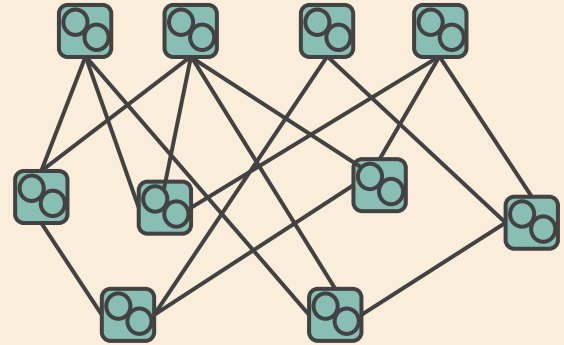


Switch
without
servers

Clos vs Expanders



Clos



Expanders

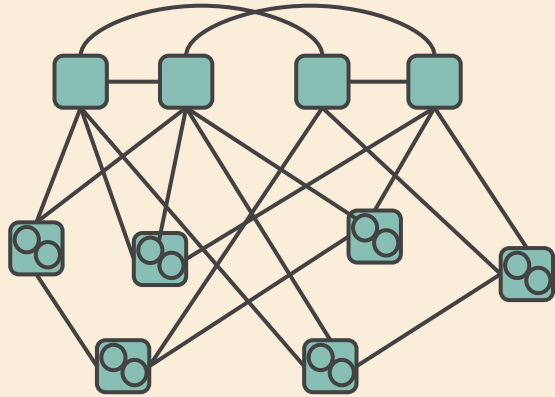


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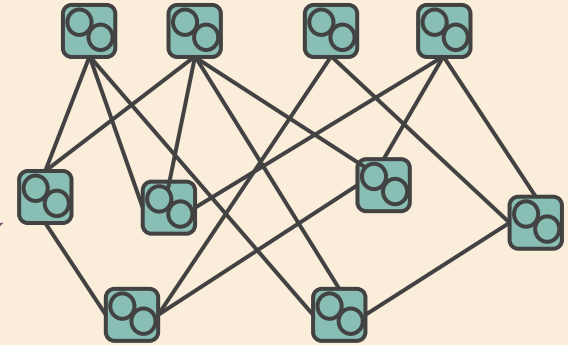
Switch
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Clos vs Expanders



Clos

2 servers



Expanders

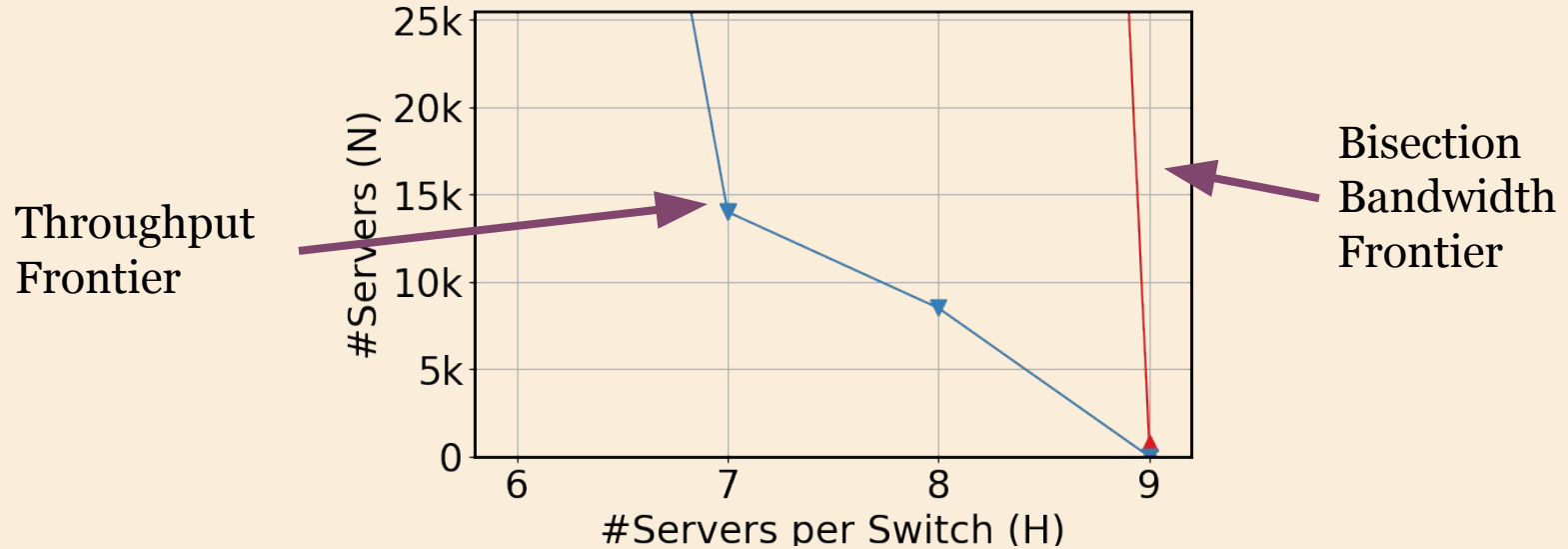


Switch
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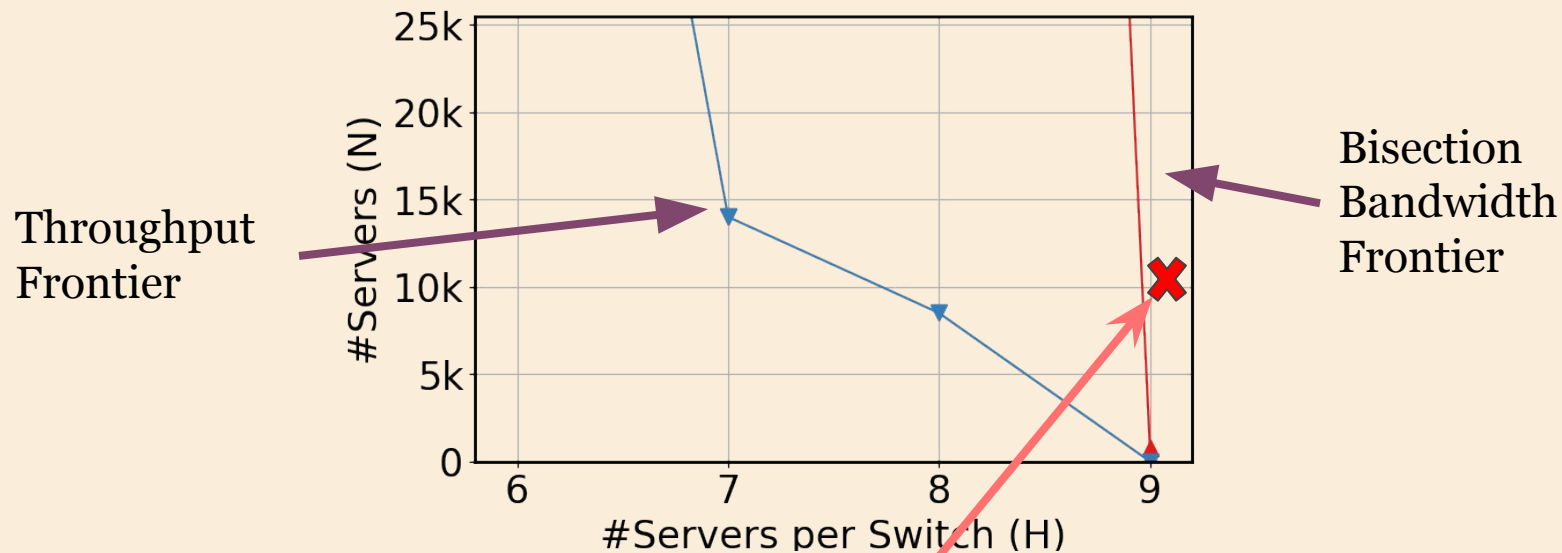


Switch
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Scaling Limitations: Frontier Curve

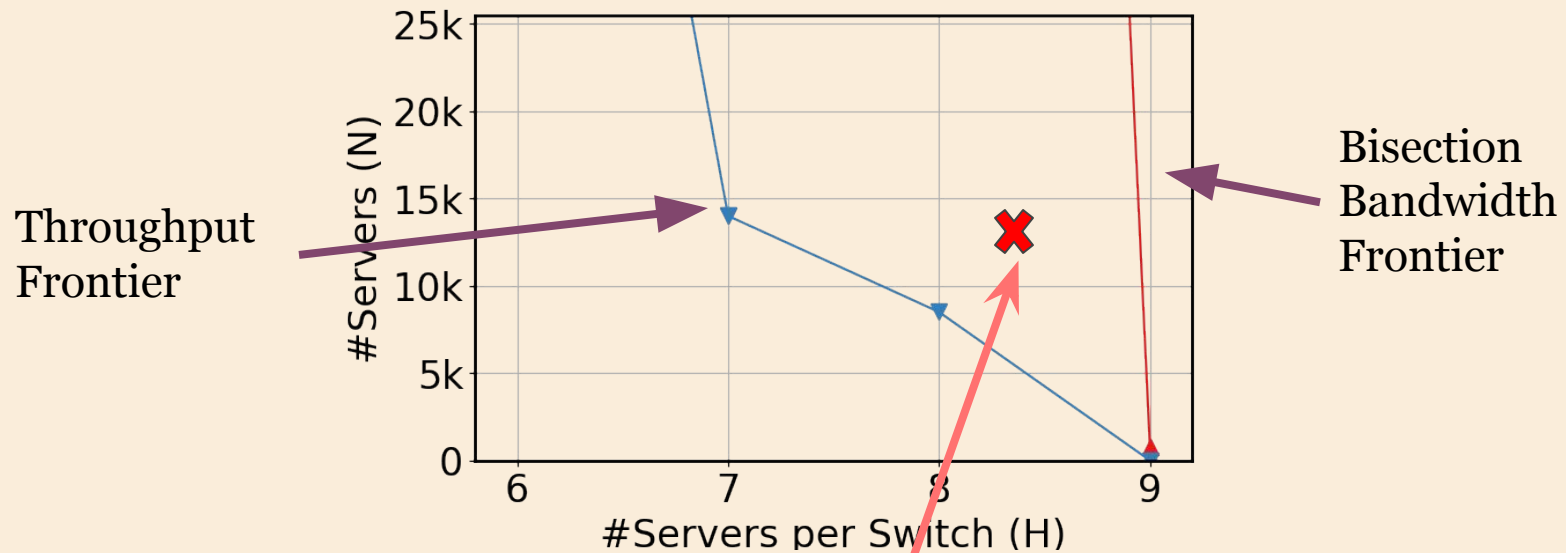


Scaling Limitations: Frontier Curve



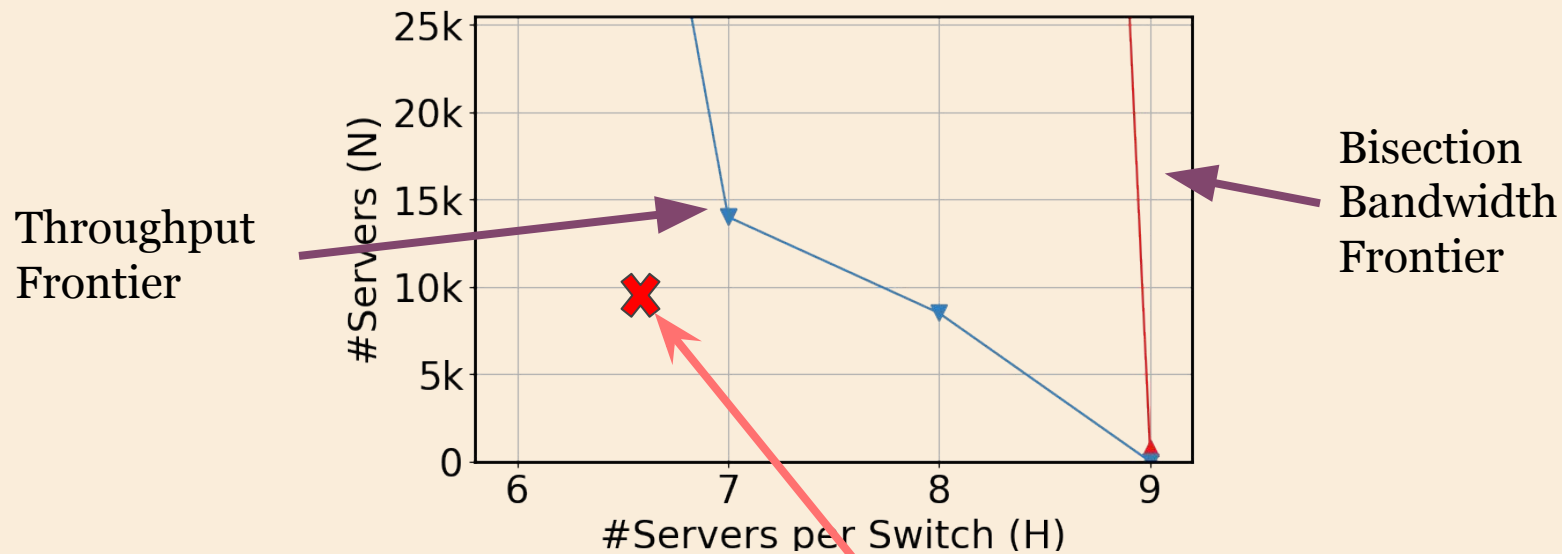
Not Full Bisection Bandwidth
Not Full Throughput

Scaling Limitations: Frontier Curve



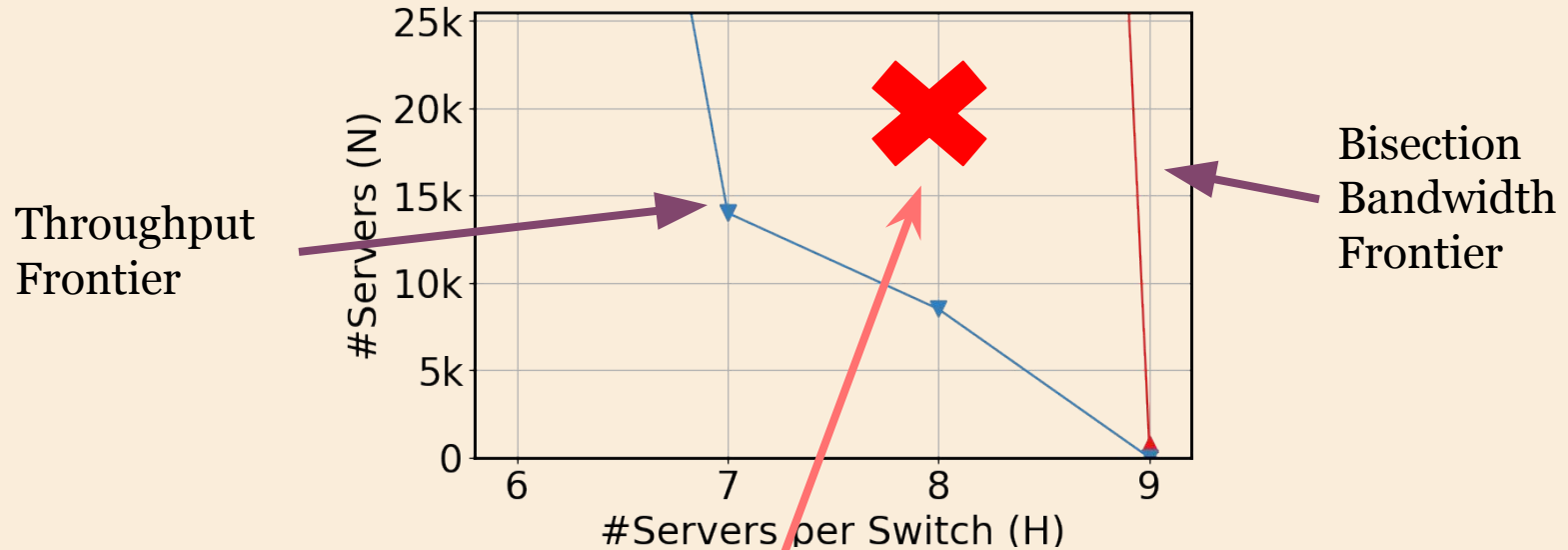
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Scaling Limitations: Frontier Curve



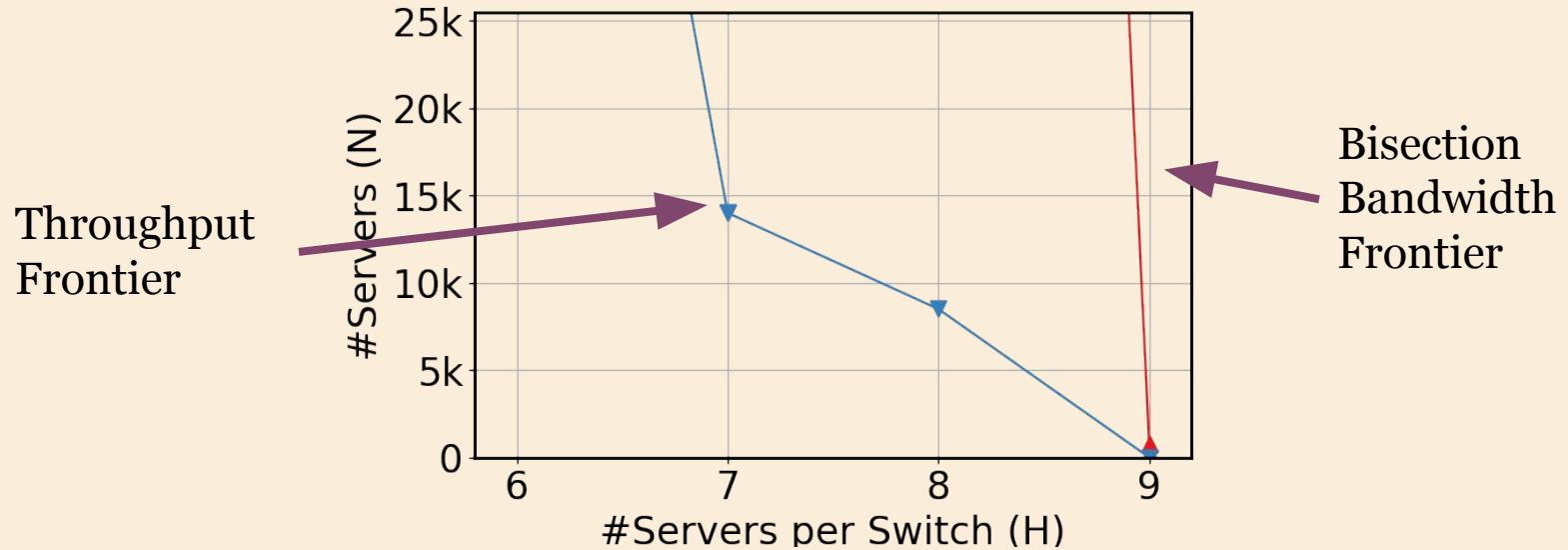
Full Bisection Bandwidth
Full Throughput

Scaling Limitations: Frontier Curve



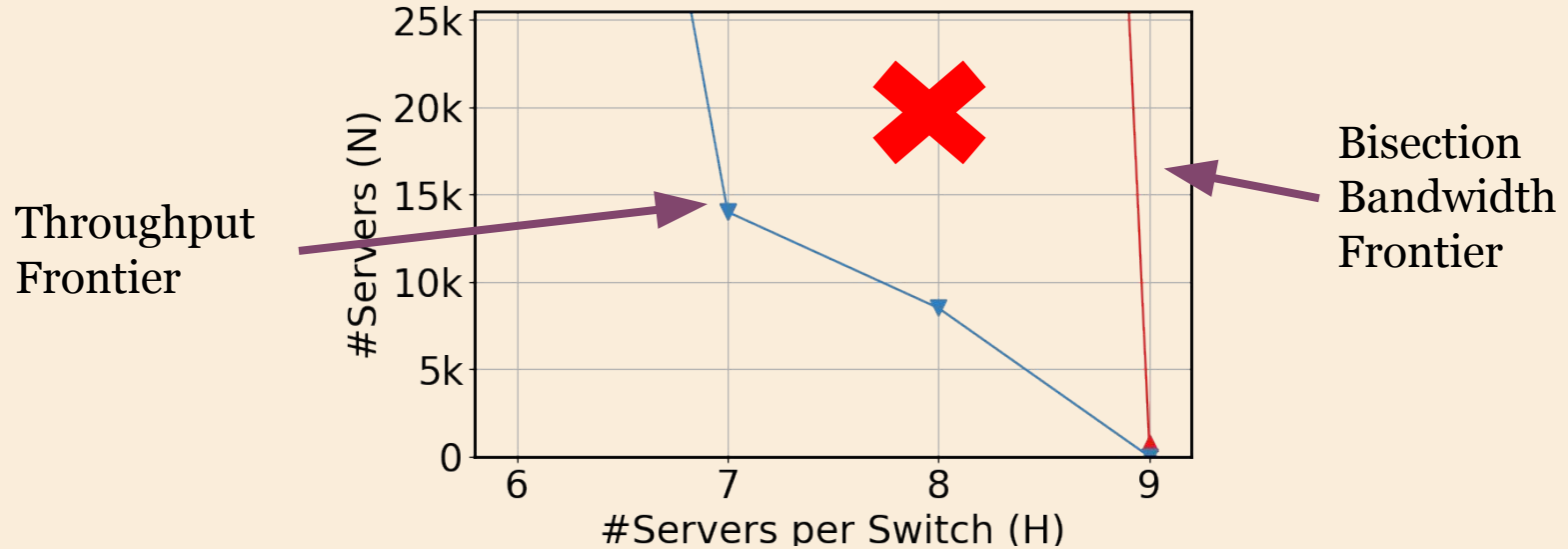
Full bisection bandwidth expanders may not be non-blocking
(not so for Clos)

Scaling Limitations: Frontier Curve



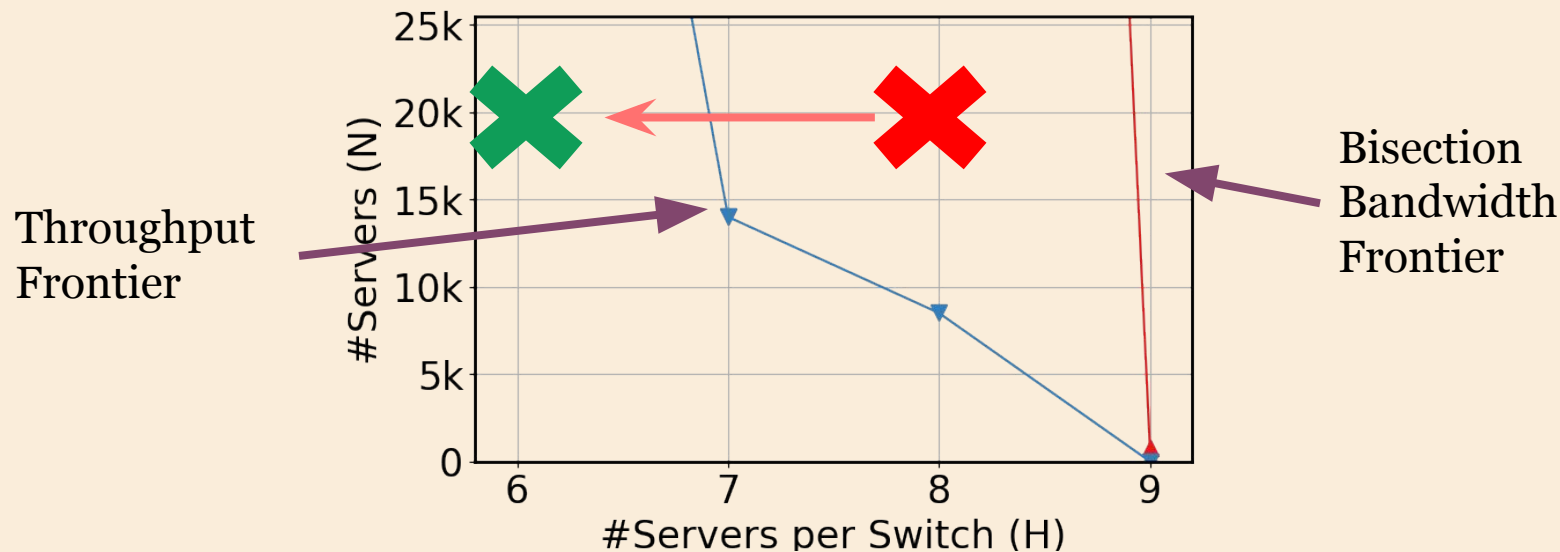
A designer may need to pick topology parameters carefully: even a small-scale expander may not be non-blocking

Scaling Limitations: Frontier Curve



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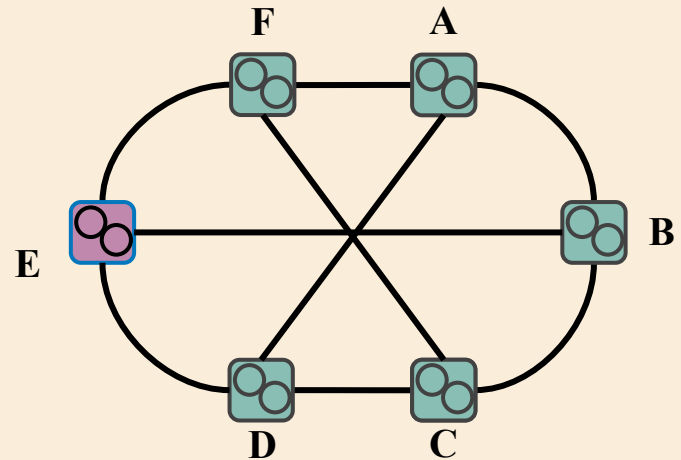
Scaling Limitations: Frontier Curve



A designer may need to pick topology parameters carefully: even a small-scale expander may not be non-blocking

Why Expanders have scaling limitations?

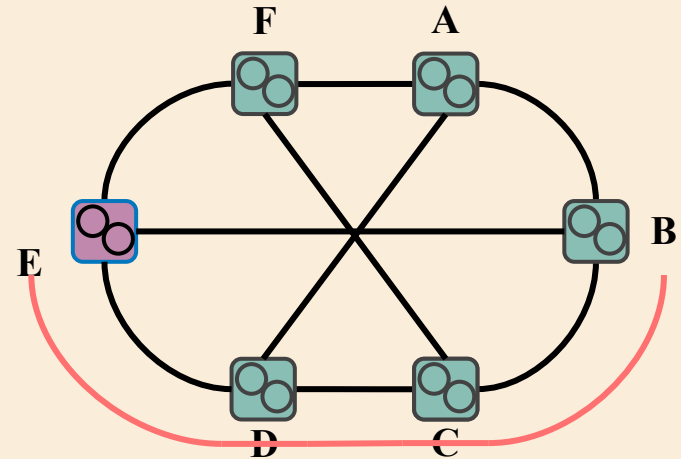
Two types of traffic in datacenter: Transit Traffic, Traffic originated/destined to connected server



Why Expanders have scaling limitations?

Two types of traffic in datacenter: Transit Traffic, Traffic originated/destined to connected server

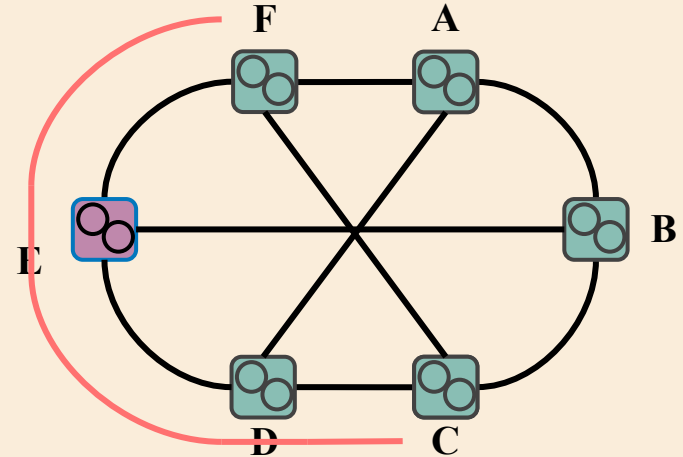
Traffic from/to the servers



Why Expanders have scaling limitations?

Two types of traffic in datacenter: Transit Traffic, Traffic originated/destined to connected server

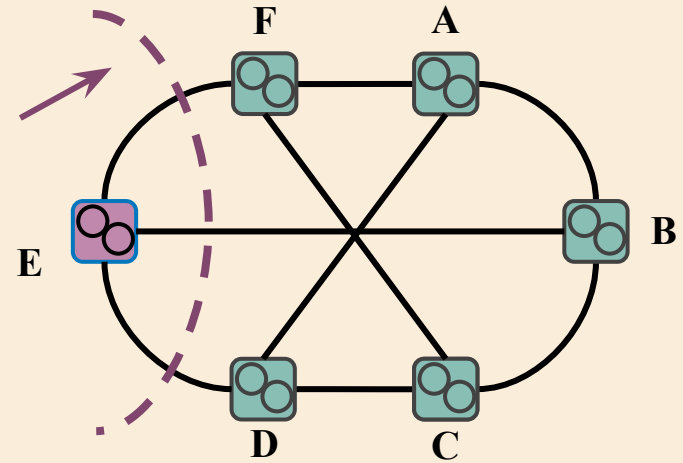
Transit Traffic



Why Expanders have scaling limitations?

Each switch has limited up-facing capacity.

Each Switch has 3 up-facing capacity

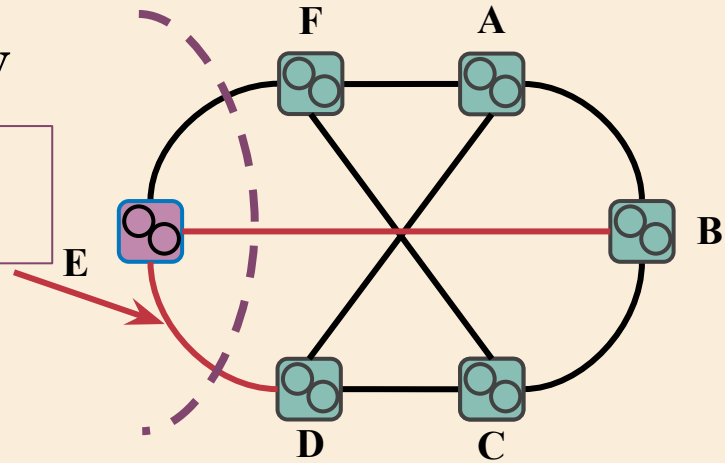


Why Expanders have scaling limitations?

In Expander, each switch has a fixed number of servers

Each Switch has 3 up-facing capacity

Each Switch connected to 2 Servers



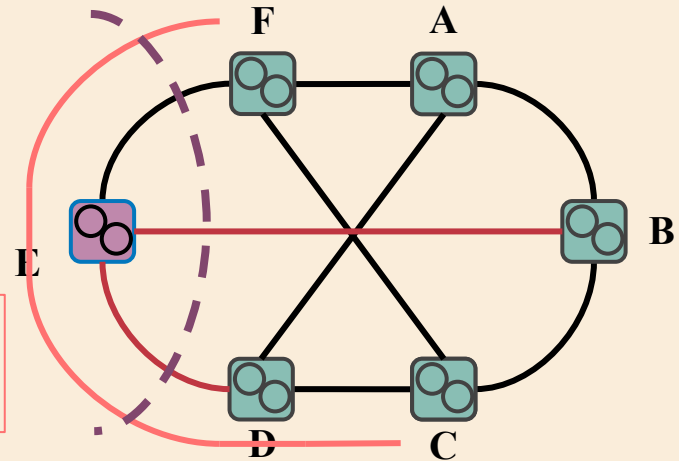
Why Expanders have scaling limitations?

In Expanders, each switch has limited capacity to handle transit traffic.

Each Switch has 3 up-facing capacity

Each Switch connected to 2 Servers

1 up-facing capacity left for transit traffic



Why Expanders have scaling limitations?

In Expanders, each switch handles both transit traffic and the traffic from/to their servers.



In Expander, number of servers per switch should be reduced so that each switch has more capacity left for transit traffic.

Conclusion

1

A full bisection bandwidth Expander may not have full throughput.

2

Cost, manageability, and failure resilience comparisons affected significantly when throughput is used at large-scale.

3

An accurate upper bound for throughput of Expanders and Clos topologies that scales well.

Future Work

- Practical routing evaluation
- Parallel Throughput upper bound computation
- Further Improvement of accuracy

Thank you!

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